Opening remarks

So far

- Basic blocks.
- Control Flow Graphs.
- Dominators, Loops
- Liveness analysis
- Register allocation (linear scan, Kempe, spilling)
- Optimizations in the basic block.
- Peephole optimizations

Announcements:

Assignment 6 is due in ten days.

Today

Runtime management - Procedure calling



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Parameter passing

Call-by-value

- store values, not addresses
- never restore on return
- arrays, structures, strings are a problem

Call-by-reference

- pass address
- access to formal is indirect reference to actual

Call-by-value-result

- store values, not addresses
- always restore on return
- arrays, structures, strings are a problem



Runtime management

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Parameter passing - varargs

What about variable length argument lists?

- if caller knows that callee expects a variable number
 - o caller can pass number as 0th parameter
 - 2 callee can find the number directly
- if caller doesn't know anything about it
 - callee must be able to determine number
 - first parameter must be closest to FP

Consider printf:

- number of parameters determined by the format string
- it assumes the numbers match



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MIPS procedure call convention

Registers:

Number	Name	Usage
0	zero	Constant 0
1	at	Reserved for assembler
2, 3	v0, v1	Expression evaluation, scalar function results
4–7	a0-a3	first 4 scalar arguments
8–15	t0t7	Temporaries, caller-saved; caller must save to preserve across calls
16–23	s0-s7	Callee-saved; must be preserved across calls
24, 25	t8, t9	Temporaries, caller-saved; caller must save to preserve across calls
26, 27	k0, k1	Reserved for OS kernel
28	gp	Pointer to global area
29	sp	Stack pointer
30	s8 (fp)	Callee-saved; must be preserved across calls
31	ra	Expression evaluation, pass return address in calls

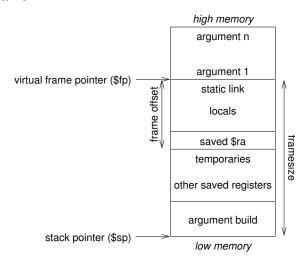
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MIPS procedure call convention

The stack frame



The "locals" can be accessed by a callee.

MIPS procedure call convention

Philosophy:

Use full, general calling sequence only when necessary; omit portions of it where possible (e.g., avoid using fp register whenever possible)

Classify routines as:

- non-leaf routines: routines that call other routines
- leaf routines: routines that do not themselves call other routines
 - leaf routines that require stack storage for locals
 - leaf routines that do not require stack storage for locals



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MIPS procedure call convention

Pre-call:

- Pass arguments: use registers \$a0 ... \$a3; remaining arguments are pushed on the stack along with save space for \$a0 ... \$a3
- Save caller-saved registers if necessary
- Execute a jal instruction: jumps to target address (callee's first instruction), saves return address in register \$ra



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MIPS procedure call convention

Prologue:

- Leaf procedures that use the stack and non-leaf procedures:
 - Allocate all stack space needed by routine:
 - local variables
 - saved registers
 - sufficient space for arguments to routines called by this routine

```
subu $sp,framesize
```

2 Save registers (\$ra, etc.):

```
sw $31, framesize+frameoffset ($sp)
sw $17, framesize+frameoffset-4 ($sp)
sw $16, framesize+frameoffset-8 ($sp)
where framesize and frameoffset (usually negative) are
compile-time constants
```

Emit code for routine



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Closing remarks

What did we do today?

- Runtime management
- Parameter passing



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MIPS procedure call convention

Epilogue:

- Opy return values into result registers (if not already there)
- Restore saved registers

```
lw reg, framesize+frameoffset-N($sp)
```

Get return address

lw \$31, framesize+frameoffset(\$sp)

Clean up stack

addu \$sp, framesize

Return

j \$31



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