

# CS6848 - Principles of Programming Languages

## Principles of Programming Languages

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## Recall

$e ::= x \mid \lambda x.e \mid e_1 e_2 \mid c \mid succ\ e$   
 $x \in \text{Identifier (infinite set of variables)}$   
 $c \in \text{Integer}$   
 $v ::= c \mid \lambda x.e$   
 $t ::= \text{Int} \mid t- > t$



## Recap

- Structural subtyping
- Unification algorithm



## Extending a language

- Extend the language grammar that will lead to new terms.
- Extend the allowed values.
- Extend the types.
- New operational semantics.
- New typing rules.



## Pairs

- Expressions

$$e ::= \dots | (e_1, e_2) | e.1 | e.2$$

- Values

$$v ::= \dots | (v_1, v_2)$$

- Types

$$t ::= \dots | t_1 \times t_2$$



## Properties of pairs

- The components are evaluated left to right.
- The pair must be fully evaluated to get the components.
- A pair that is passed as an argument will be fully evaluated, before the function starts executing (in call by value semantics).



## New operational semantics

- First element:

$$(Pair \beta 1) \quad (v_1, v_2).1 \rightarrow v_1$$

- Second element:

$$(Pair \beta 2) \quad (v_1, v_2).2 \rightarrow v_2$$

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$$[Projection 1] \quad e \rightarrow e' \implies e.1 \rightarrow e'.1$$

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$$[Projection 2] \quad e \rightarrow e' \implies e.2 \rightarrow e'.2$$

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$$[Pair Evaluation 1] \quad e_1 \rightarrow e'_1 \implies (e_1, e_2) \rightarrow (e'_1, e_2)$$

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$$[Pair Evaluation 2] \quad e_2 \rightarrow e'_2 \implies (v_1, e_2) \rightarrow (v_1, e'_2)$$



## Tuples

- Expressions

$$e ::= \dots | (e_i^{i \in 1..n}) | e.i$$

- Values

$$v ::= \dots | (v_i^{i \in 1..n})$$

- Types

$$t ::= \dots | (t_i^{i \in 1..n})$$



- Element  $j$ :

$$(\beta)(v_i^{i \in 1..n}).j \rightarrow v_j$$

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$$[\textit{Projection 1}]e \rightarrow e'.e.i \rightarrow e'.i$$

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$$[\textit{Tuple Evaluation}]e_j \rightarrow e'_j(v_i^{i \in 1..j-1}, e_j, e_k^{k \in j+1..n}) \rightarrow (v_i^{i \in 1..j-1}, e'_j, e_k^{k \in j+1..n})$$



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$$[\textit{Tuple}]A \vdash \forall i e_i : t_i A \vdash (l_i = e_i^{i \in 1..n}) : (l_i : t_i^{i \in 1..n})$$

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$$[\textit{Projection}]A \vdash e : (l_i : t_i^{i \in 1..n}) A \vdash e.l_j : t_j$$

