Assignment #2 CS3300

1. [5] **Activation records**: Draw the runtime activation records at the entry and exit to each function. Pay special attention to the following fields of the activation records: actual parameters, return value, control link, and local variables.

```
foo(){ // Returns the number of primes in x.
  int x[] = {2, 3, 4, 5, 6}, len = 5;
  return Eratosthenes(x, len, 0);
}
int Eratosthenes(int x[], int len, int index){
  int i, v = x[index];
  for (i=index+1; i<len; i+= v) x[i] = -1;
  for (i=index+1; i< len; i++)</pre>
    if (x[i] != -1) return Eratosthenes (x, len, i);
  return countNonNegative (x, len);
}
int countNonNegative(int x[], int len){
  int i, count=0;
  for (i=0;i<len;++i) if (x[i] > 0) count ++;
  return count:
}
```

2. [5] **Code generation**: Write an SDT to generate IR in three-address code (similar to the one discussed in the class) for the following grammar. Briefly explain about the attributes you use.

```
\begin{array}{l} \mathsf{P} \rightarrow \mathsf{S} \\ \mathsf{S} \rightarrow \mathsf{SwitchStmt} \; \mathsf{S} \; \mid \; \mathsf{Assignment}; \; \mathsf{S} \; \mid \; \epsilon \\ \mathsf{SwitchStmt} \rightarrow \mathsf{switch} \; ( \; \mathrm{Id} \; ) \; \{ \; \mathsf{CaseBlocks} \; \} \\ \mathsf{CaseBlocks} \rightarrow \mathsf{case} \; \mathsf{Lit} \; : \; \mathsf{Stmt2}; \; \mathsf{CaseBlocks} \; \mid \; \mathsf{default}: \; \mathsf{Stmt2} \\ \mathsf{Stmt2} \rightarrow \mathsf{Assignment} \; \mid \; \mathsf{Break} \\ \mathsf{Assignment} \rightarrow \mathsf{x} \; = \; \mathsf{E} \\ \mathsf{E} \; : \; \mathsf{RelEx} \; \mid \; \mathsf{AddEx} \; \mid \; \mathsf{Id} \\ \mathsf{RelEx} \rightarrow \; \mathsf{E} \; < \; \mathsf{E} \\ \mathsf{AddEx} \rightarrow \; \mathsf{E} \; + \; \mathsf{E} \\ \mathsf{Break} \; \rightarrow \; \mathsf{break} \end{array}
```

3. [5] Flow-graph: Generate the three-address code and draw the flow graph.

```
void sort (int A[], int n){
  for (int i=0;i<n-1;++i)
    for (int j=i+1;j<n; ++j)
        if (A[i] > A[j]){ tmp = A[i]; A[i] = A[j]; A[j] = tmp; }
}
```

- 4. [5] **Register Allocation**: Prove that the live variable analysis discussed in the class computes the liveness information for each variable, conservatively. At runtime, if the value of a variable v is live at a program point L (that is used at a later point of L), then the computed In(L) set includes v.
- 5. [5] **Personal!** Based on the portion covered after Quiz 1 in CS3300, make an interesting question and attempt an answer for the same. Credit will be given based on the creativity, ingenuity, and coolness of the question and the answer.