# CS3300 Quiz 2: Oct 10, 2023. (QP Code: A) <br> Maximum marks $=30$, Time: 45 minutes, Closed Book, Closed Neighbor 

Name: $\qquad$ Roll: $\qquad$
Read the instructions and questions carefully. You can use the given booklet for rough work and stating any reasonable assumptions you make. But write the answers in the QP itself - marks will be given based on the answers in the QP.

- MCQ and True/False questions: Each incorrect answer will lead to a deduction of 0.5 marks.


## Section 1. MCQ, 2 marks each

1. 

Consider the action sequence shown to process one of the production rules to translate the boolean predicate $B$.

B1.true = B2.true = B.true
Assume the used attributes have the same meaning dis-
B $\rightarrow$ B1 || B2
B1.false $=$ B2.false $=$ B.false
cussed in the class.
B.code $=$ B1.code || B2.code

Which of the following IS true:
(a) During execution of the generated code, the code for B1 will always jump to the correct statement.
(b) During execution of the generated code, the code for B 2 will never be executed.
(c) During execution of the generated code, the code for B will lead to an infinite loop.
(d) None of the above.
2. Consider a C array X, declared as "int X[10] [20]". Students of CS3300 came out with four different translations for the statement $\mathrm{a}=\mathrm{X}[\mathrm{i}][\mathrm{j}]$, as shown below:

| $t 1=j * 40$ | $t 1=i * 80$ | $t 1=i * 40$ | $t 1=j * 80$ |
| :--- | :--- | :--- | :--- |
| $t 2=i * 4$ | $t 2=j * 4$ | $t 2=j * 4$ | $t 2=i * 4$ |
| $t 3=t 1+t 2$ | $t 3=t 1+t 2$ | $t 3=t 1+t 2$ | $t 3=t 1+t 2$ |
| $a=X[t 3]$ | $a=X[t 3]$ | $a=X[t 3]$ | $a=X[t 3]$ |
| (A) | (B) | (C) | (D) |

t2=i*4
$\mathrm{t} 3=\mathrm{t} 1+\mathrm{t} 2 \quad \mathrm{t} 3=\mathrm{t} 1+\mathrm{t} 2$
(A)
(B)

Assuming that the size of an integer is 4 bytes and the array is allocated in column-major form, which is the correct three-address IR code.
(a) A
(b) B
(c) C
(d) D
3. Which of the following is true?
(a) Every $\operatorname{LR}(\mathrm{k})$ grammar can be rewritten as a $\mathrm{LR}(1)$ grammar?
(b) $\operatorname{LALR}(1)$ grammar may have conflicts
(c) For arbitrary values of $k_{1}$ and $k_{2}$, in general, the expressive power of $\operatorname{LR}\left(k_{1}\right)$ is not comparable to that LL $\left(k_{2}\right)$.
(d) JavaCC uses LR(k) parsing.

## Section 2. Fill in the blank, 4 marks each

Consider the C code shown in RHS:
The number of nodes and edges in the CFG for its three-address-code are $\qquad$ and $\qquad$ or ( $1=1 ; 1<=n ;++i$ )
for ( $\mathrm{j}=1 ; \mathrm{j}<=\mathrm{n} ;++\mathrm{j}$ ) $\{$
C [i] $=0$;

1. respectively. [ $2+2$ marks].

Hint: Don't forget to include the entry and exit blocks. for ( $k=1 ; k<=n ;++k)\{$
$\mathrm{C}[\mathrm{i}]+=\mathrm{A}[\mathrm{i}] * \mathrm{~B}[\mathrm{i}] ;$
\}
\}

Consider the LR(1) item sets shown on the right.
(a) The number of shift entries in the $\mathrm{LR}(1)$ parsing table $=$ $\qquad$ . [1 mark]
(b) The number of reduce en-
2. tries in the LR(1) parsing table = $\qquad$ [1 mark]
(c) The number of rows (excluding the header row) in the $\operatorname{LALR}(1)$ table would be $\qquad$ . [1 mark]

(d) State True or False: the grammar is $\operatorname{LALR}(1)$ $\ldots$. [1 mark]

Consider the two C codes shown on the right
(a) The number of jump instructions (count both conditional and unconditional) in the generated code, for the code A and code B are $\qquad$ and $\qquad$ _, respectively. [2 marks]

```
i = 0;
while (i < n){
while (i < n)
        i++;
}
```

i $=0$;
if (i < n) \{
do \{
A[i] = i;
i++;
\} while ( $\mathrm{i}<\mathrm{n}$ );
\}
(A)
(B)
(b) During execution, the number of executed jump instructions (count both conditional and unconditional), for the code $A$ and code $B$ loop are $\qquad$ and $\qquad$ , respectively. State your answer in terms of n . [2 marks]
Consider the CFG shown on the right.
4.
(a) The number of production rules in the augmented grammar $=$ $\qquad$ [1 mark].

$$
\begin{aligned}
& S \rightarrow a S b \\
& S \rightarrow A \\
& A \rightarrow a A \\
& A \rightarrow \epsilon
\end{aligned}
$$

(b) The number of $\operatorname{LR}(1)$ item-sets for this gram$\operatorname{mar}=$ $\qquad$ [3 marks].
Section 3. True or False Answers, 1 mark each
_ The size of the AST is typically smaller than the syntax tree.
In any CFG, the number of join nodes will always match that of the branch nodes.
___ Given an LL(1) grammar, top-down parsers cannot perform on-the-fly evaluation of attributes in Sattributed grammars.
_ A CFG is a rooted undirected graph.
$\qquad$ An Inherited attribute of a node $N$ is defined only in terms of attributes at the children of $N$ and at $N$.
$\qquad$ $\operatorname{LR}(1)$ is strictly more expressive than $\operatorname{LL}(k)$ (irrespective of the value of $k$ ).
Given an attribute dependence graph, with no cycles, there is a unique order for evaluating the rules.
L-attributed grammars are a superset of S-attributed grammars.

