

Final Exam

CS6013

Maximum marks = 50, Time: 3hrs

23-Nov-2012

1. [10] **Dominance Frontiers**

Define dominance frontier and write an algorithm to compute the Dominance Frontiers for all the nodes in the CFG.

Input: Set of Nodes N , Set of Edges E , root node r

Output: $\forall v \in N, DF(v)$

Assume that you are given the following maps:

- **Succ** and **Pred** functions, that return the successor and predecessors respectively.
- **DOM**, **SDOM**, **IDOM** functions return the dominators, strict dominators, and immediate dominators for a given node.

What is the complexity of the algorithm?

[*Bonus* 5] Can you implement a faster algorithm?

2. [15] **Register allocation**

a) Briefly describe the iterated register coalescing algorithm.

b) Intra-procedural register allocation suffers from many drawbacks: the values present in different registers before a call are copied to argument registers / swap area; the caller and callee save registers have to appropriately saved. Extend the iterated register coalescing algorithm to design an inter-procedural register allocation scheme that avoids these copy related issues. Assume that standard intra-procedural liveness information is available. Assume that the globals are live throughout the program.

```
b=1;                               r1 = 1;
c=2;                               r2 = 2;
a = foo(b, c).                      call foo
print a;                            print r3;
p = c;                               r1 = r2;
q = 3;                               r2 = 3;
x = foo (p, q);                      →   call foo
print x;                            print r3
int foo(int f, int g){              foo:
    k = f op g                       r3 = r1 op r2
    return k;
}
```

3. [15] **Copy propagation**

Given an assignment $x := y$, the *copy propagation* optimization, replaces each later use of x with y , as long as there is no redefinition of x or y in between. Write an iterative algorithm to do intra-procedural copy propagation. Apply your algorithm on the following code:

```
    c = a + b
    d = c
    e = d * d
L1: f = a + c
    g = e
    a = g + d
    if (a < c) {
        h = g + 1
        L2: b = g * a
            if (h < f) goto L1
    }
    else {
        f = d - g
        if (f > a) goto L2
    }
    else
        c = 2
    }
```

4. [10] **Dependence analysis**

Answer the below mentioned queries for the given sample code:

```
for (i = 1; i <= n; ++i) do
    for (j = n; j >= 1; --j) do
        for (k = 1; k <= n+1; ++k) do
            S1: A[i,j,k] = A[i,j-1,k-1] + A[i-1,j,k]
            S2: B[i,j-1,k] = A[i,j-1,k-1] * 2.0
            S3: A[i,j,k+1] = B[i,j,k] + 1.0
        endfor
    endfor
endfor
```

- Draw the iteration space for the loop nest.
- Draw the execution order relationships between the three labeled statements.
- Write the dependence relations (flow, anti and output) between the three labeled statements.
- Restate the dependence relations in terms of distance vectors, direction vectors and dependence vectors.
- For the different references to A and B, use the GCD test to check if/when there exists any same iteration or different iteration dependence.