

# CS3300

## Quiz 2

Dept of CSE, IIT Madras

Total marks = 24

Time = 50 min

15 Oct 2019

**Read the instructions and questions carefully.** You can make any reasonable assumptions that you think are necessary; but state them clearly. There are total four questions ( $8 + 8 + 4 + 4 = 24$  marks + 2 bonus marks). You will need approximately 15 minutes for answering an 8 marks question (plan your time accordingly). For questions with sub-parts, the division for the sub-parts are given in square brackets.

You will get an answer sheet with 8 pages (if you get a answer booklet with fewer pages then ask for a replacement). Leave the first page empty and start from Page#2. Start each question on a new page. Think about the question before you start writing and write briefly. **For any question, the answer (including the answers for all the sub-parts) should NOT cross more than two pages.** If the answer for any question is spanning more than two pages, we will strictly ignore the spill-over text. If you scratch/cross some part of the answer, you can use space from the next page. You mostly would NOT need any additional sheets.

1. [8+1] **Syntax Directed Translation:** Define inherited and synthesized attributes [2]. Given a production  $A \rightarrow X_1 X_2 \dots X_n$ , state the order in which the inherited and synthesized attributes of  $A$ ,  $X_1$ ,  $X_2$ ,  $\dots$   $X_n$  may be evaluated in an L-attributed grammar [1]. Give an example grammar and semantic rules (not discussed in the slides) illustrating both inherited and synthesized attributes [Bonus 1].  
Write a grammar to recognize positive decimal integer literals in C [1]. Use this grammar to write an SDT scheme to check if the number is divisible by six. [4]. You may not use division (or repeated subtraction for that matter).
2. [8] **Type checking and IR Generation:**  
Consider a subset of MiniJava with no loops and no conditional statements. To translate this subset, give a minimal list of instructions of 3-address-codes (TAC), with one line description for each [3]. For each such TAC instruction, give an example MiniJava code snippet (we don't require the complete program) explaining where we need that TAC instruction [2]. Assume that the IR need not have a statement for variable declarations. If you want you can also use 'bOp' as the generic binary operator.  
Give rules on how you will type-check a program written in your TAC [3].
3. [4+1] **Liveness Analysis and register allocation**  
Using a simple example code, explain **def** and **use** [1]. Say we have a program with no branch instructions then the liveness analysis algorithm can be simplified significantly. Write the simplified algorithm [3] and discuss its complexity [Bonus 1].
4. [4] **Answer True or False**
  - (a) In an L-attributed grammar some attributes can be both synthesized and inherited.
  - (b) Every compiler for a given language, for a given pass (for example, type checking) must use the same data-structure for the symbol table.
  - (c) If MiniJava had method overloading but no method overriding, then the call resolution can happen at compile time.
  - (d) LL parsers push semantic actions into the parse stack like grammar symbols.
  - (e) In the procedure linkages, the prologue and epilogue code is exactly opposite of each other.
  - (f) Given a program, during the discussed iterative liveness analysis algorithm, the **use** and **def** information may change, but only monotonically.
  - (g) Given the set of live-in variables of a node, we cannot precisely compute its live-out variables using the set of variables used/defined at that node.
  - (h) Liveness analysis cannot be performed as a forward analysis.