

## Hoare Triples

- Syntax:  $\{P\} S \{Q\}$
- Intended meaning: If statement  $S$  is executed starting in an initial state satisfying  $P$  and it terminates, the resulting final state will satisfy  $Q$

## Exercise

- Write a specification
- ... for a program  $S$  that swaps the values of two variables  $A$  and  $B$

## Exercise

- Is the following valid?

```

{n > 0}
sum := 0;
i := 0;
while (i < n) {
  sum := sum + 1;
  i := i + 1;
}
{ sum == n*(n+1)/2 }

```

## Program Correctness

- Partial correctness vs. total correctness
- Proving program termination usually harder and requires different techniques

## Exercise

- What will really happen if we run this program (in practice today)?

```

{ n > 0 }
sum := 0;
i := 0;
while (i < n) {
  sum := sum + 1;
  i := i + 1;
}
{ sum == n*(n+1)/2 }

```

## Integers in practice

- Integer variables & arithmetic
  - Classical (standard) theory
  - Reality: 32-bit and 64-bit integers
  - Bit-vector arithmetic theory
  - Modular arithmetic

## Arrays

- Specify that an array  $A$  is sorted
- How can we model arrays?
- The logic we have considered so far is untyped. We can extend this to a typed version (called many-sorted logic)
  - “sort”  $\approx$  “type”
- We will stick to plain untyped logic here

## The Array Datatype

- Function symbols used to model arrays
  - **length** (unary)
    - **length(A)**
      - We can use  $A.length$  as syntactic sugar
  - **lookup** (binary)
    - **lookup(A,i)**
      - We can use  $A[i]$  as syntactic sugar
  - **update** (ternary)
    - **update(A,i,v)**
      - We can use  $A[i \leftarrow v]$  as syntactic sugar
      - same as array  $A$  except at index  $i$  where it contains the value  $v$
  - **isArray** (unary)
    - type information

## Exercise

- Specify that an array  $A$  is sorted
- Specify that a program  $S$  sorts an input array  $A$
- What are appropriate axioms for the array datatype?

## Pointers & Heap

- Memory can be modeled as one (or more) giant array
- Naïve modeling may not be sufficient
  - Due to incompleteness
- Separation Logic
- Pointer analysis
- TVLA
  - An abstract-interpretation based approach to verification in the presence of dynamic structures

## Hoare Logic

- Inference rules for proving  $\vdash \{P\} S \{Q\}$
- Combines
  - Rules for standard mathematical logic with
  - Rules for reasoning about programming language constructs

Hoare Logic

## Basic Constructs

- Assignment statement:  $x := e$
- Statement sequencing:  $S_1; S_2$
- Conditional: *if* ( $E$ ) *then*  $S_1$  *else*  $S_2$
- Iteration: *while* ( $e$ ) *do*  $S$

## Statement Sequencing

$$\frac{?}{\vdash \{P\} S_1; S_2 \{Q\}}$$

$$\frac{\vdash \{P\} S_1 \{R\}, \quad \vdash \{R\} S_2 \{Q\}}{\vdash \{P\} S_1; S_2 \{Q\}}$$

## Conditional Statement

$$\frac{?}{\vdash \{P\} \text{if } (E) \text{ then } S_1 \text{ else } S_2 \{Q\}}$$

$$\frac{\vdash \{P \wedge E\} S_1 \{Q\}, \quad \vdash \{P \wedge \neg E\} S_2 \{Q\}}{\vdash \{P\} \text{if } (E) \text{ then } S_1 \text{ else } S_2 \{Q\}}$$

## Iteration

$$\frac{?}{\vdash \{P\} \text{while } (E) \text{ do } S \{Q\}}$$

$$\frac{\vdash \{P \wedge E\} S \{P\}, \quad \vdash (P \wedge \neg E) \Rightarrow Q}{\vdash \{P\} \text{while } (E) \text{ do } S \{Q\}}$$

## Assignment

$$\frac{?}{\vdash \{P\} x := E \{Q\}}$$

$$\frac{}{\vdash \{Q[x \rightarrow E]\} x := E \{Q\}}$$

## Exercise

- Try out the assignment rule
- $\{?\} y := 2 * x \{y > 10\}$
- $\{?\} x := x + 1 \{x > 10\}$

## Concurrency

- Separation Logic