Pointer Analysis

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Dead Code Elimination

```
a = s1.arr;
b = s2.ptr;
q = &a[ii];
p = &b[jj];
if(p == q) {
     x = 10;

y = 100;
y = 30;
```

To check the condition, we need to test if

- p == q• a + ii * typesize == b + jj * typesize
- s1.arr + ii * typesize == s2.ptr + jj * typesize

This needs to be tested statically

Outline

- Introduction
- Pointer analysis as a DFA problem
- · Design decisions
- Andersen's analysis, Steensgaard's analysis
- Pointer analysis as a graph problem
 - Optimizations
- · Applications
- Parallelization
 - Constraint based
 - Replication based
 - Graph rewrite rules

Common Subexpression Elimination

```
q = s1.arr;
p = s1.ptr;
if(p + i == q + j) \{
    x = 10;

y = 100;
} else {
    x = 20;
    y = 30;
```

To identify if the expression is common

- p + $i\dot{i}$ == q + $j\dot{j}$ s1.arr + $i\dot{i}$ * typesize_ $i\dot{i}$ == s1.ptr + $j\dot{j}$ * typesize_ $j\dot{j}$

This needs to be computed statically

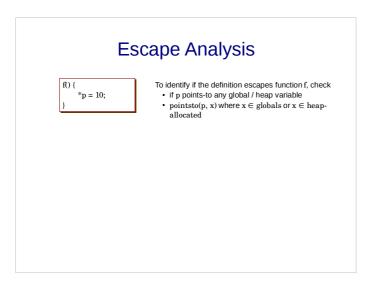
Applications

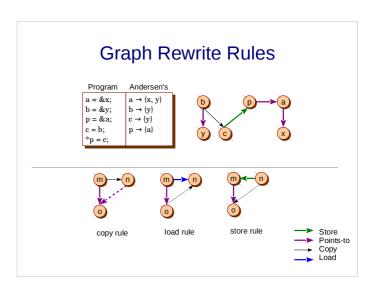
- Dead-code elimination
- · Common subexpression elimination
- Parallelization
- · Escape analysis

Parallelization

To identify if the functions are parallelizable, check if

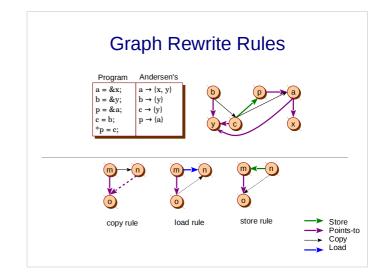
• !alias(*p, *q)





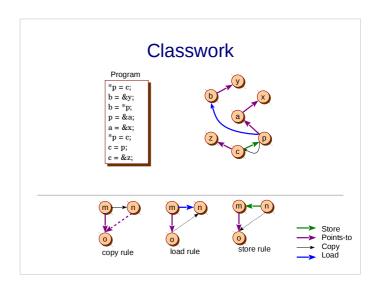
Parallel Pointer Analysis

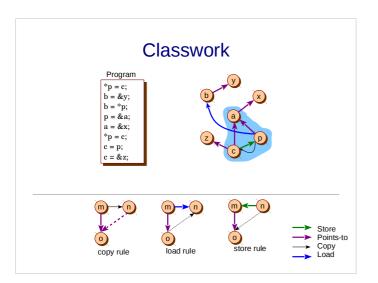
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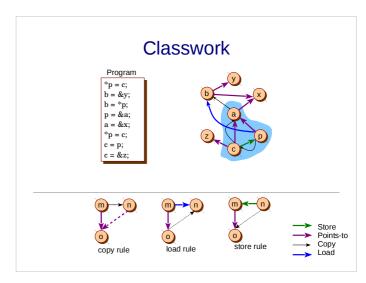


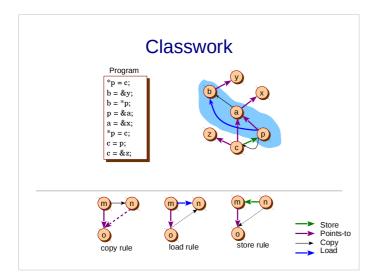
Pointer Analysis as Graph Rewrite Rules

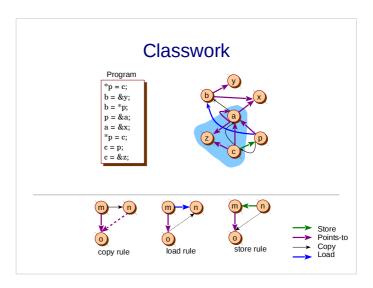
- Initially: Constraint-based: pointers and associated points-to sets
- Later: Graph problem: pointers as nodes, subset relation forms edges, points-to set with each node
- Now: Graph rewrite rules: variables as nodes, all relations form edges, points-to set defined using edges

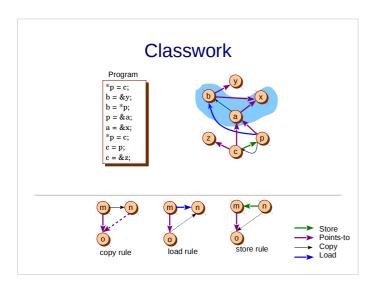


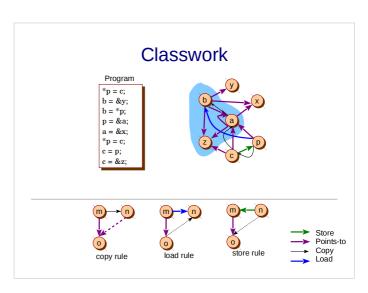


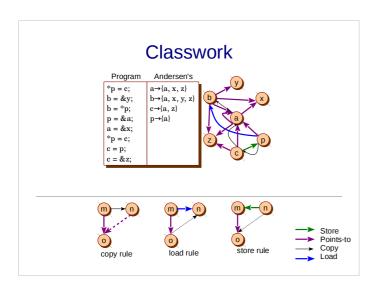












Parallel Graph Rewrite Rules

- Open: How to order rule evaluation?
- *Open:* How to combine rules for better efficiency?

