

Shape Analysis

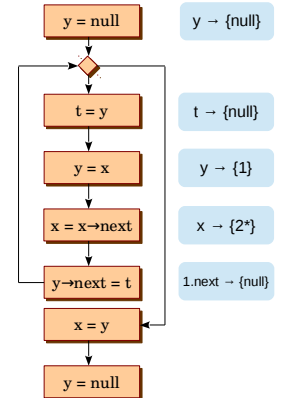
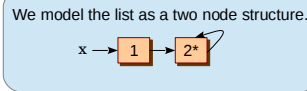
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CS6843 Program Analysis
IIT Madras
Jan 2016

Limitations of Pointer Analysis

```
listReverse(List x) {
  assert("x is an acyclic singly linked list");

  for (y = null; x;) {
    t = y;
    y = x;
    x = x->next;
    y->next = t;
  }
  x = y;
  t = null;
  y = null;
}
```



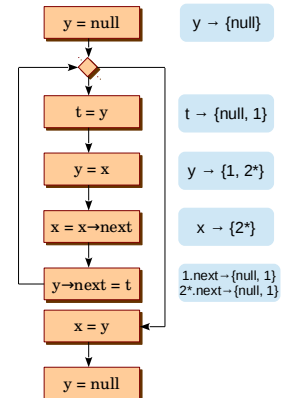
Outline

- Limitations of pointer analysis
- Identify lists
- Identify trees, DAGs, cyclic graphs
- Identifying rotations
- List reversal and other transformations

Limitations of Pointer Analysis

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}
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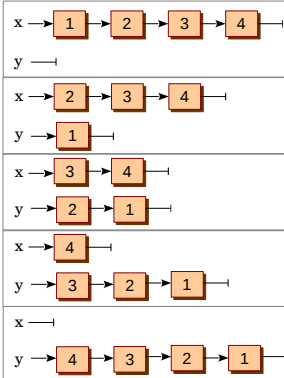


Limitations of Pointer Analysis

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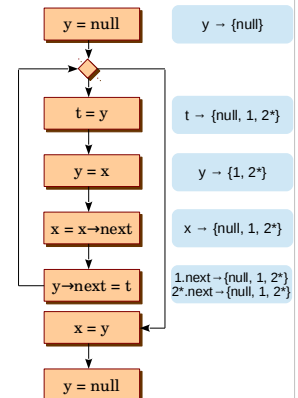
We want to check if x points to a singly linked list at the end of listReverse. That is, $x \rightarrow \{4\}$, $4.next \rightarrow \{3\}$, ..., $1.next \rightarrow \{null\}$



Limitations of Pointer Analysis

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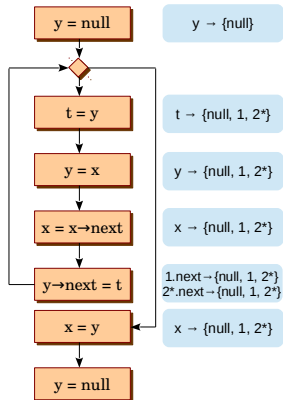
  for (y = null; x;) {
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Limitations of Pointer Analysis

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listReverse(List x) {
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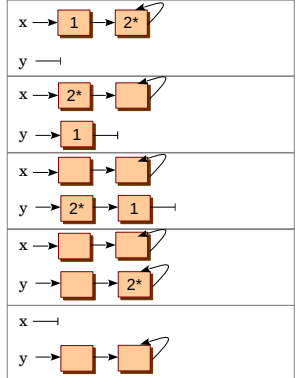
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Shape Analysis

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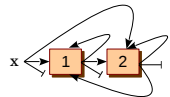


Maintain additional information with 2^* that it is acyclic. Use the fact that node removal maintains acyclicity.

Limitations of Pointer Analysis

```
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    x = x->next;
    y->next = t;
  }
  x = y;
  t = null;
  y = null;
}
```



Annotations for the diagram:

- $y \rightarrow \{\text{null}\}$
- $t \rightarrow \{\text{null}, 1, 2^*\}$
- $y \rightarrow \{\text{null}, 1, 2^*\}$
- $x \rightarrow \{\text{null}, 1, 2^*\}$
- $1.\text{next} \rightarrow \{\text{null}, 1, 2^*\}$
 $2^*.\text{next} \rightarrow \{\text{null}, 1, 2^*\}$
- $x \rightarrow \{\text{null}, 1, 2^*\}$

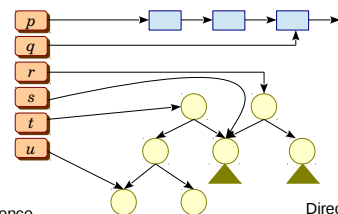
Tree, DAG, Cycle?

- Maintains three data structures:
 - Interference matrix: encodes common reachability
 - Direction matrix: encodes direct reachability
 - Shape
- Performs iterative data-flow analysis to update D, I and shape information

Shape Analysis

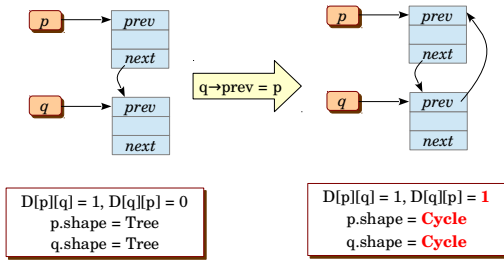
- Identify structural / topological properties of a data structure under manipulation.
- Usually categorized as slist, tree, DAG or cycle.
- Precision reduces along slist \rightarrow tree \rightarrow DAG \rightarrow cycle.

Tree, DAG, Cycle?



Interference							Direction						
	p	q	r	s	t	u		p	q	r	s	t	u
p	1	1	0	0	0	0	p	1	1	0	0	0	0
q		1	0	0	0	0	q	0	1	0	0	0	0
r			1	1	1	0	r	0	0	1	1	0	0
s				1	1	0	s	0	0	0	1	0	0
t					1	1	t	0	0	0	1	1	1
u						1	u	0	0	0	0	0	1

Shape Estimation



Inference Rules

$p = malloc(...)$	D_kill and I_kill sets same as for allocation statement.
$p = q$	$D_gen = \{D[s][p] \mid D[s][q] \text{ and } s \neq p\} \cup$ $\{D[p][s] \mid D[q][s] \text{ and } s \neq p\} \cup$ $\{D[p][p] \mid D[q][q]\}$
$p = q \rightarrow f$	
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = null$	$I_gen = \{I[p][s] \mid I[q][s] \text{ and } s \neq p\} \cup$ $\{I[p][p] \mid I[q][q]\}$
$p \rightarrow f = q$	
$p \rightarrow f = null$	$p.shape = q.shape$

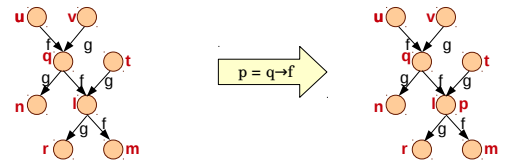
The implementation should create new D/I matrices from their current copies. In-situ update would lead to unsound or imprecise analysis.

Inference Rules

$p = malloc(...)$	
$p = q$	
$p = q \rightarrow f$	
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = null$	
$p \rightarrow f = q$	
$p \rightarrow f = null$	

Inference Rules

$p = malloc(...)$	D_kill and I_kill sets same as for allocation statement.
$p = q$	$D_gen = \{D[s][p] \mid I[s][q] \text{ and } s \neq p\} \cup$ $\{D[p][s] \mid D[q][s] \text{ and } s \neq p \text{ and } s \neq q\} \cup$ $\{D[p][q] \mid q.shape == Cycle\} \cup$ $\{D[p][p] \mid D[q][q]\}$
$p = q \rightarrow f$	
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = null$	
$p \rightarrow f = q$	$I_gen = \{I[p][s] \mid I[q][s] \text{ and } s \neq p\} \cup$ $\{I[p][p] \mid I[q][q]\}$
$p \rightarrow f = null$	$p.shape = q.shape$



Inference Rules

$p = malloc(...)$	$D_kill = \{D[p][s] \mid D[p][s] == 1\} \cup \{D[s][p] \mid D[s][p] == 1\}$ $I_kill = \{I[p][s] \mid I[p][s] == 1\}$
$p = q$	$D_gen = \{D[p][p]\}$
$p = q \rightarrow f$	$I_gen = \{I[p][p]\}$
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = null$	$p.shape = Tree$
$p \rightarrow f = q$	
$p \rightarrow f = null$	

Inference Rules

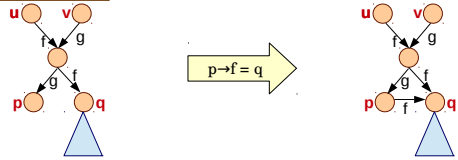
$p = malloc(...)$	Processing is the same as for $p = q$ statement. This means the analysis loses field-sensitivity.
$p = q$	
$p = q \rightarrow f$	
$p = \&(q \rightarrow f)$	A former work from IITK (Dasgupta, Karkare, Reddy) addresses this issue.
$p = q \text{ op } k$	
$p = null$	
$p \rightarrow f = q$	
$p \rightarrow f = null$	

Inference Rules

<p>$p = \text{malloc}(\dots)$</p> <p>$p = q$</p> <p>$p = q \rightarrow f$</p> <p>$p = \&(q \rightarrow f)$</p> <p>$p = q \text{ op } k$</p> <p>$p = \text{null}$</p> <p>$p \rightarrow f = q$</p> <p>$p \rightarrow f = \text{null}$</p>	<p>D_kill and I_kill sets same as for allocation statement.</p> <p>$D_gen = \{ \}$</p> <p>$I_gen = \{ \}$</p> <p>$p.\text{shape} = \text{Tree}$</p>
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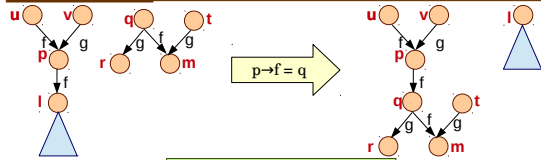
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---	---



Inference Rules

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---	---



Can you improve precision?

Inference Rules

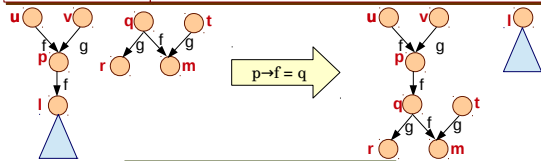
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---	---

	Tree	DAG	Cycle
Tree	Tree	DAG	Cycle
DAG	DAG	DAG	Cycle
Cycle	Cycle	Cycle	Cycle

max(shape1, shape2)

Inference Rules

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For max() consider $D[v][r] == 1$.

Inference Rules

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---	--

Example

```
listReverse(List x) {
  assert("x is an acyclic singly linked list");

  for (y = null; x;) {
    t = y;
    y = x;
    x = x->next;
    y->next = t;
  }
  x = y;
  t = null;
  y = null;
}
```

Interference

	x	y	t
x	0	0	0
y		0	0
t			0

Direction

	x	y	t
x	1	0	0
y	0	0	0
t	0	0	0

Shape

x	tree
y	tree
t	tree

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We need to assume a finite representation for the data structure.

Interference

	x	y	t
x	1	0	0
y		0	0
t			0

Direction

	x	y	t
x	1	0	0
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Shape

x	tree
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Interference

	x	y	t
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```

Since we do not model fields, we can't say that x is unreachable from y.

Interference

	x	y	t
x	1	0	0
y		0	0
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x	y	t
x	1	1
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t		

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Direction

x	y	t	
x	1	0	0
y	1	1	1
t	1	1	1

Shape

x	tree
y	cycle
t	cycle

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Classwork

```
p = malloc(10);
p->f1 = null;
q = p->f2;
q = &(r->f2);
q->f2 = p;
```

Example

```
listReverse(List x) {
  assert("x is an acyclic singly linked list");

  for (y = null; x;) {
    t = y;
    y = x;
    x = x->next;
    y->next = t;
  }
  x = y;
  t = null;
  y = null;
}
```

Interference

x	y	t
x	1	1
y		1
t		

Direction

x	y	t
x	1	0
y	1	1
t	1	1

Shape

x	tree
y	cycle
t	cycle

Improvements

- Field-sensitivity
- Heap modeling
- Path-sensitivity

Example

```
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Interference

x	y	t
x	1	1
y		1
t		

Direction

x	y	t
x	1	1
y	1	1
t	1	1

Shape

x	cycle
y	cycle
t	cycle

Summary

- Shape analysis helps several transforms.
- Existing techniques often trade off precision for efficiency.
- We are still far away from a precise and scalable analysis.

