### **Dynamic Analysis**

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## Outline

- Applications of dynamic analysis
  - Limitations of static analysis
  - Trade-offs
- Profiling techniques
- Finding invariants
  - Equality
  - Affine
- Dynamic type inferencing

## Applications

- Bug finding (testing)
- Data race detection
- Identifying security vulnerabilities
- Improved precision of static analysis
- Input-dependent analysis

### Limitations of Static Analysis

- Reduced precision: Over-approximations
- Cannot perform input-dependent analysis

### Static versus Dynamic

Sound

Incomplete

- Imprecise
- Input-oblivious

- Precise
- Input-dependent

- Choosing between static and dynamic analysis often requires a trade-off between soundness and precision.
- Current trend is to combine the two techniques to get better precision at improved scalability.

## Profiling

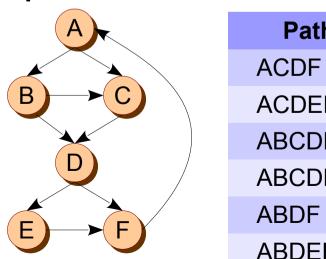
- Profiling is a method of collecting information of interest during program execution.
- The information is often useful to find hot-spots in the program.
- Examples
  - Number of times an instruction is executed
  - Number of page faults
  - Number of cache hits
  - Total memory used

# Profiling

- Intrusive: inserts instructions in the program (source, IR, assembly) statically, which get executed at runtime
  - File log
  - Memory locations pointed to by a pointer
  - Execution time of a function
- Non-intrusive: the program is unaltered; uses external means to profile
  - Hardware counters
  - Program execution time

### Path Profiling

- Consider a program with an entry node and an exit node. There are several execution paths (traces) that the program takes from entry to exit.
- The task is to find the frequency of execution of each path.



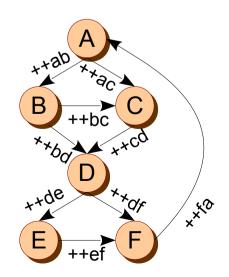
Path	Frequency
ACDF	90
ACDEF	60
ABCDF	0
ABCDEF	100
ABDF	20
ABDEF	0

### Path Profiling

- Naïve path profiling is expensive: instrumenting each path may lead to exponential blow up in computation and storage
- This can lead to unacceptable program slowdown

### Edge Profiling

- Path profile is approximated as an edge profile
- The frequency of each edge is calculated which is used to find the path frequency



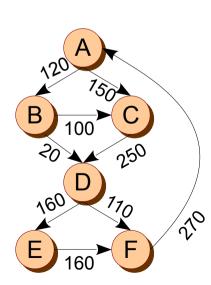
### Edge Profiling

- Path profile is approximated as an edge profile
- The frequency of each edge is calculated which is used to find the path frequency

Path	Frequency	Choose the minimum
ACDF	110	edge-frequency in the path
ACDEF	150 🧹	
ABCDF	100	
ABCDEF	100	
ABDF	20	
ABDEF	20	11
	ACDF ACDEF ABCDF ABCDEF ABDF	ACDF110ACDEF150ABCDF100ABCDEF100ABDF20

### Path vs. Edge Profiling

- Path profile is approximated as an edge profile
- The frequency of each edge is calculated which is used to find the path frequency
  - Can this instrumentation be optimized?Can we have better precision?



Path	Path Frequency (actual)	Path Frequency (estimated)
ACDF	90	110
ACDEF	60	150
ABCDF	0	100
ABCDEF	100	100
ABDF	20	20
ABDEF	0	20

### Efficient Edge Profiling

- Observation: We do not need to instrument every edge.
- How to find a minimal, low-cost set of edges to instrument?
- Use a spanning tree (instrument non-st edges):
  - reduced instrumentation along paths,
  - not all edges carry instrumentation

Classwork: Find counts for uninstrumented edges.

A	Path	Frequency
***	$C\toD$	ac + bc
++bc	$D\toF$	ac + bc + bd - de
	$E \to F$	de
e e e e e e e e e e e e e e e e e e e	$A \rightarrow B$	bc + bd
► F	$F\toA$	ac + bc + bd

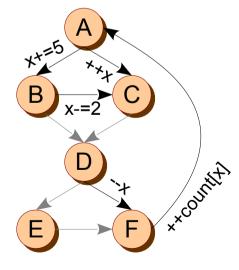
## Edge Profiling

- Edge profile may not always be a good indicator of a path profile
- Efficient edge profiling requires a unique variable along each instrumented edge (non-spanning tree edge)

120 A	Path	Frequency	Actual Freq.	Actual Freq. 2
150	ACDF	110	90	110
	ACDEF	150	60	40
250	ABCDF	100	0	0
160 770	ABCDEF	100	100	100
	ABDF	20	20	0
	ABDEF	20	0	20

#### But path profiling is expensive

- Since index variable across all paths
- Path linearization: Unique (and consecutive) path numbering, which enables indexing
- Most hardware support registers, fast increment and indexing

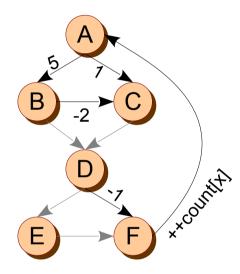


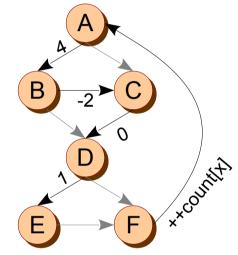
Path	x
ACDF	0
ACDEF	1
ABCDF	2
ABCDEF	3
ABDF	4
ABDEF	5

**Classwork**: Prove that such a path numbering is unique.

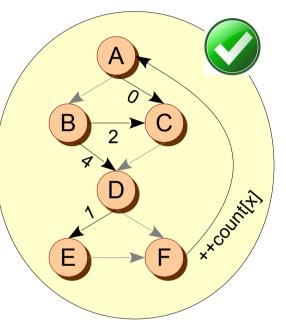
Check the value of x for each path.

Path numbering is not unique





Path	x
ACDF	0
ACDEF	1
ABCDF	2
ABCDEF	3
ABDF	4
ABDEF	5



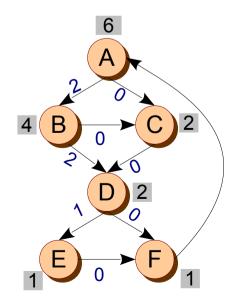
In all the above cases, the path numbering is the same, number of instrumented edges (5) is the same

So, which instrumentation should we choose?

- 1. Assign integer values to edges such that no two paths compute the same path-sum.
- 2. Use a spanning tree to select edges to instrument and compute the appropriate increment.
- 3. Select appropriate instrumentation.
- 4. After collecting the run-time profile, derive the execution paths.

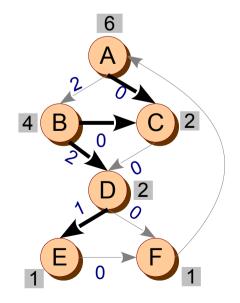
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NumPaths(node) = 0 NumPaths(leaf) = 1 In reverse topological order For each edge  $v \rightarrow w$  {  $Val(v \rightarrow w) = NumPaths(v)$  NumPaths(v) += NumPaths(w)}



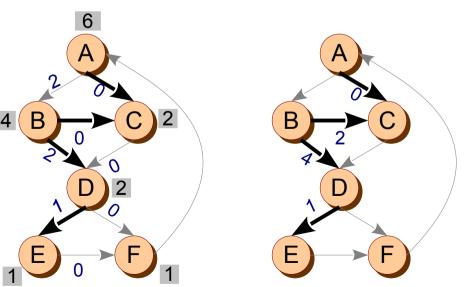
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- Find a spanning tree.
- Find chord (non-ST) edges.
- For each chord, find fundamental cycle.



- 1. Assign integer values to edges such that no two paths compute the same path-sum.
- 2. Use a spanning tree to select edges to instrument and compute the appropriate increment.
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Chord AC: cycle ACDF : 0 Chord BC: cycle ABCDF : 2 Chord BD: cycle ABDF : 4 Chord DE: cycle DEF : 1



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*Prelude:* Allocate and initialize the array of counters

# *Postlude:* Write the array to permanent storage

Main:

- Initialize path register r in the entry vertex
- Increment path memory counter in the exit vertex
- Optimizations

- 1. Assign integer values to edges such that no two paths compute the same path-sum.
- 2. Use a spanning tree to select edges to instrument and compute the appropriate increment.
- 3. Select appropriate instrumentation.
- 4. After collecting the run-time profile, derive the execution paths.

#### Path Regeneration Path id $\rightarrow$ Path mapping?

A	Path	id
2 0	ACDF	0
B → C \	ACDEF	1
2 0	ABCDF	2
	ABCDEF	3
	ABDF	4
	ABDEF	5

### Classwork

• Find the instrumentation for the following CFG

