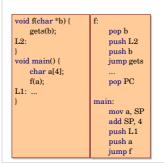
# **Security Analysis**

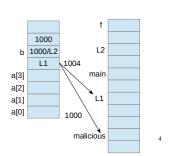
Rupesh Nasre.

CS6843 Program Analysis IIT Madras Jan 2016

## Stack Smashing

 How can a malicious code be executed by exploiting buffer overrun vulnerability?





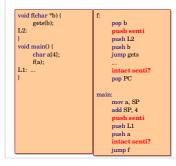
#### **Outline**

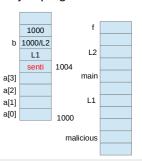
- · Introduction and applications
- · Buffer overrun vulnerability

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# To Avoid Stack Smashing

- Insert a sentinel near the return address.
- · Check if it is intact before jumping.





#### Introduction

- · Security in a broad sense.
  - Effects: crash, non-termination, wrong output, unintended actions
  - Causes: dangling pointers, buffer overruns, null pointer dereference, wrong opcode, arbitrary data-change
- C programs are more susceptible to buffer overflow attacks.
- C allows direct pointer manipulation since space and performance are primary concerns not security.
- Standard library contains functions that are unsafe if not used carefully (e.g., gets, strcpy, strcat). Does strncpy solve the problem?

### To Avoid Stack Smashing

- Insert sentinel / canary
- Check addresses / bounds explicitly (Java)
- · Wrap system calls with security checks

Dynamic techniques
 Runtime overhead
 Program is terminated

- When the code segment is writable, it is more vulnerable to attacks (self-modifying code, W^X).
- · What does the following program do?

char\*f="char\*f=%c%s%c;main(){printf(f,34,f,34,10);}%c";main(){printf(f,34,f,34,10);}

## Notes on Stack Smashing

- Using canary for stack smashing detection?
  - Canary is a bird used in coal-mines to detect toxic gases (humans follow the caged birds)
  - Researchers have validated its performance impact to be minimal
  - Randomizing canary improves odds
  - Does not guarantee protection
- · How about heap smashing?
  - Heap usually doesn't contain return addresses
  - But then, we have function pointers

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## Specifying Pre and Post-conditions

```
char *strcpy(char *s1, char *s2)
```

```
/* @requires maxDef(s1) >= maxDef(s2) */
/* @ensures maxUse(s1) == maxUse(s2)
```

and result == s1 \*/;

void \*malloc(size\_t size)

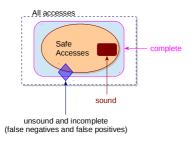
/\* @ensures maxDef(result) == size

or result == null \*/;

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#### Static Buffer Overrun Detection

 A good example of static analysis that can be incomplete as well as unsound.



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#### **Inferring Constraints**

- From the for-loops init, bound and change
  - Difficult for general loops such as while
- From the array declarations and malloc statements
- · From conditional checks in the code
- Small number of heuristics often cover large part of the program.
- Once the constraints are identified, these are checked against the user annotations.

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# Using Pre and Post-conditions

- Annotations define properties
  - minDef, maxDef, minUse, maxUse

e.g., minDef(buff) = 0, maxUse(buff) = N / 2

- notNull, null, restrict

e.g., notNull(ptr), restrict(ptr)

- Homework: Write an example program using restrict which enables an optimized code.
- Initially we would assume that these annotations are user-provided. Later, we will try to auto-infer them.

# **Inferring Constraints**

- In absence of annotations, simply generating all possible constraints is expensive.
- In the past, researchers have tried flowinsensitive constraints.
- Auto-inference is feasible when loop-bounds do not depend on array values.
  - while  $(a[i] != '\0')$  versus while (i < n)

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## Precision vs. Efficiency

- Precision requires interprocedural analysis in the above example

  (recell Analysis Dimensions)
- (recall Analysis Dimensions).Domain knowledge about N may help in filtering out false positives.

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# Vulnerability Analysis in Polyhedral Model

- How do you model inequalities?
- · What are the constants?
- What do you get after solving the system?

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# Stack Smashing in gcc



#### **Tools**

#### 3. BOON

- Array out of bound check for C
- Flow-insensitive, intra-procedural, pointer-insensitive

#### 2. CQual

- Annotation-based
- Uses type qualifiers to propagate taint annotation
- Detects format string vulnerability by type checking

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# Vulnerability Analysis as a DFA

- · Data-flow facts
- · Statements of interest
- · Analysis direction
- · Meet operator

Classwork

#### **Tools**

#### 1. xq++

- Template-driven compiler extension
- Finds kernel vulnerabilities
- Tracks kernel data originated in untrusted source, memory leaks, deadlock situations

#### 0. Eau Claire

- Theorem-prover based (specification-checker)
- Finds buffer overruns, file access races, format string bugs

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