

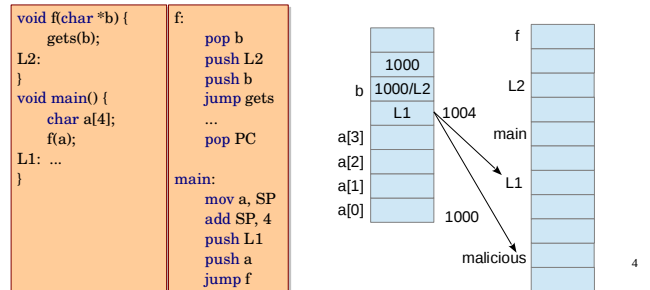
## Security Analysis

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CS6843 Program Analysis  
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## Stack Smashing

- How can a malicious code be executed by exploiting buffer overrun vulnerability?



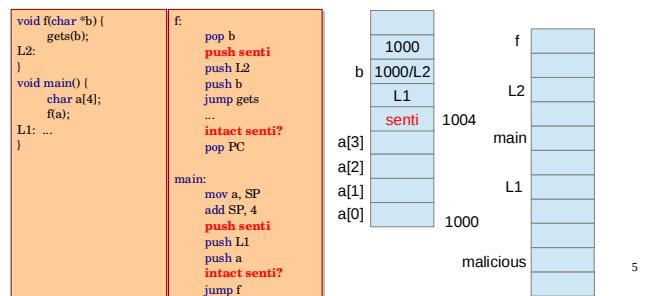
## Outline

- Introduction and applications
- Buffer overrun vulnerability

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## To Avoid Stack Smashing

- Insert a sentinel near the return address.
- Check if it is intact before jumping.



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## Introduction

- Security in a broad sense.
  - Effects: crash, non-termination, wrong output, unintended actions
  - Causes: dangling pointers, buffer overruns, null pointer dereference, wrong opcode, arbitrary data-change
- C programs are more susceptible to buffer overflow attacks.
- C allows direct pointer manipulation – since space and performance are primary concerns – not security.
- Standard library contains functions that are unsafe if not used carefully (e.g., `gets`, `strcpy`, `strcat`). Does `strncpy` solve the problem?

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## To Avoid Stack Smashing

- Insert sentinel / canary
- Check addresses / bounds explicitly (Java)
- Wrap system calls with security checks
  - Dynamic techniques
    - Runtime overhead
    - Program is terminated
- When the code segment is writable, it is more vulnerable to attacks (*self-modifying code*, *W^X*).
- What does the following program do?

```
char*f="char*=%c%s%c;main(){printf(f,34,f,34,10);}%c";main(){printf(f,34,f,34,10);}
```

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## Notes on Stack Smashing

- Using canary for stack smashing detection?
  - Canary is a bird used in coal-mines to detect toxic gases (humans follow the caged birds)
  - Researchers have validated its performance impact to be minimal
  - Randomizing canary improves odds
  - Does not *guarantee* protection
- How about heap smashing?
  - Heap usually doesn't contain return addresses
  - But then, we have function pointers

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## Specifying Pre and Post-conditions

- `char *strcpy(char *s1, char *s2)`

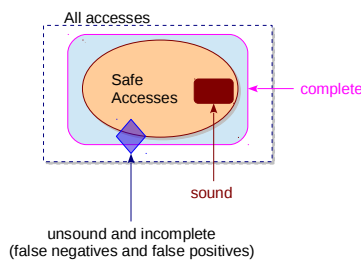
```
/* @requires maxDef(s1) >= maxDef(s2) */
/* @ensures maxUse(s1) == maxUse(s2)
   and result == s1 */;
```
- `void *malloc(size_t size)`

```
/* @ensures maxDef(result) == size
   or result == null */;
```

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## Static Buffer Overrun Detection

- A good example of static analysis that can be incomplete as well as unsound.



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## Inferring Constraints

- From the `for`-loops init, bound and change
  - Difficult for general loops such as `while`
- From the array declarations and `malloc` statements
- From conditional checks in the code
- Small number of heuristics often cover large part of the program.
- Once the constraints are identified, these are checked against the user annotations.

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## Using Pre and Post-conditions

- Annotations define properties
  - `minDef`, `maxDef`, `minUse`, `maxUse`  
e.g., `minDef(buff) = 0`, `maxUse(buff) = N / 2`
  - `notNull`, `null`, `restrict`  
e.g., `notNull(ptr)`, `restrict(ptr)`
  - **Homework:** Write an example program using `restrict` which enables an optimized code.
- Initially we would assume that these annotations are user-provided. Later, we will try to auto-infer them.

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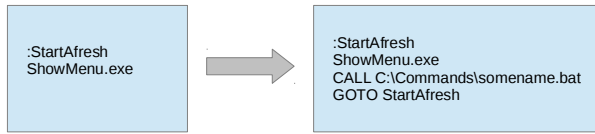
## Inferring Constraints

- In absence of annotations, simply generating all possible constraints is expensive.
- In the past, researchers have tried flow-insensitive constraints.
- Auto-inference is feasible when loop-bounds do not depend on array **values**.
  - `while (a[i] != '\0')` **versus** `while (i < n)`

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## Self-Modifying Code



Original batch file

Modified batch file

In earlier single-window DOS systems, only one window could be active, and easy inter-process communication was not well-developed.