

Roll No: _____

CS6843 Program Analysis
MidSem February 29, 2016

Duration: 60 minutes

Answer in the same sheet.

1. For the following program, draw the control-flow graph.

[3 marks]

```
1 main() {
2   int x = 0;
3   int n = 1;
4   int a = 2;
5   int i = 0;

6   if (n > 0) {
7     for (; i < n; ++i) {
8       a = i * i;
9       if (a < 100) {
10        x += a;
11      } else {
12        x -= a;
13      }
14    }
15  }
16  printf("%d\n", x);
17 }
```

2. For the above program, perform live-variables analysis. Equations are: $in(B) = use(B) \cup (out(B) - def(B))$, $out(B) = \cup in(S)$ where S is a successor of B . Simply fill-up the following table. [4 marks]

	in1	out1	in2	out2	in3	out3	outfinal
0							
1							
2							
3							
4							
5							
6							
7							

3. If your analysis is tracking one bit each for conditions $x == 0$ and $x > 2$ and $y > 4$, find the bit-values after every program statement below. Conservatively, each bit is set to 0. Your analysis does not have any other information, apart from that x and y are unsigned integers. [5 marks]

```
{ 0 0 0 }
y = 2;
{       }
x = y + 1;
{       }
x = x - y - 1;
{       }
y = y * y + 4;
{       }
x = y - 1;
{       }
```

4. For the following set of statements, perform Andersen's analysis using constraint graph and show only the final state of the graph (nodes, their attributes and edges). [5 marks]

*p = q, *a = b, c = a, p = &c, d = &a, q = *a, a = &b, *b = p, b = &p, q = &a.

5. Prof. George claims that while analyzing a program, if for every copy + load + store of the form $P = Q$, we add $Q = P$ and run Andersen's analysis, we would get the same answer as Steensgaard's analysis run on the original program. Disprove his claim with a counter-example (use as less number of statements as you can). Make sure your counter-example is reasonable. [3 marks]