OPERATING SYSTEMS CS3500

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OCT - 2022.

FILE-SYSTEM INTERFACE

Outline

- > File Concept
- Access Methods
- Disk and Directory Structure
- > Protection
- ➤ Memory-Mapped Files

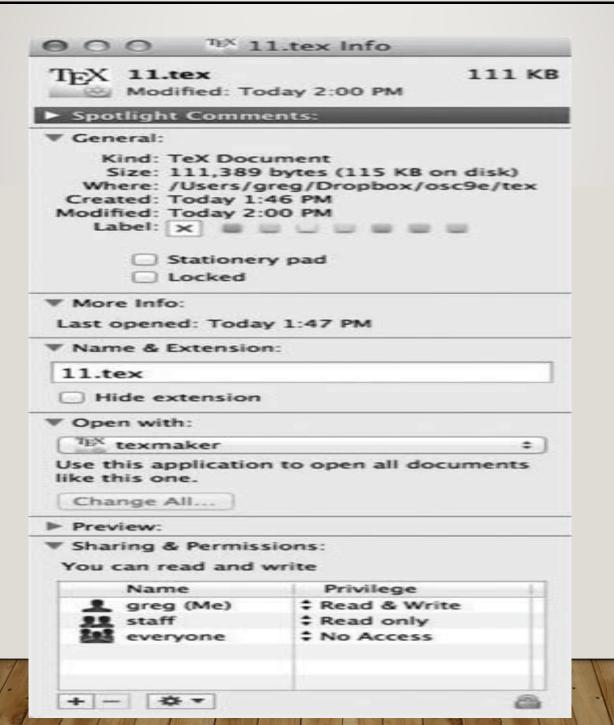
FILE CONCEPT

- A file is a named collection of related information that is recorded on secondary storage
- User's perspective: a file is the smallest allotment of logical secondary storage.
- Types:
 - Data
 - Numeric, Character, Binary
 - Program
- In general, a file is a sequence of bits, bytes, lines, or records, the meaning of which is defined by the file's creator and user
 - Many types: text file, source file, executable file

FILE ATTRIBUTES

- Name only information kept in human-readable form
- Identifier unique tag (number) identifies file within file system
- **Type** needed for systems that support different types
- Location pointer to file location on device
- Size current file size
- **Protection** controls who can do reading, writing, executing
- **Time, date, and user identification** data for protection, security, and usage monitoring
- Information about files are kept in the directory structure, which is maintained on the disk
- Many variations, including extended file attributes such as file checksum
- Information kept in the directory structure

FILE INFO WINDOW ON MAC OS X



FILE OPERATIONS

- Create
- Write at write pointer location
- Read at read pointer location
- Reposition within file seek
- Delete
- Truncate
- Open (F_i) search the directory structure on disk for entry F_i , and move the content of entry to memory
- Close (F_i) move the content of entry F_i in memory to directory structure on disk

OPEN FILES

- Several pieces of data are needed to manage open files:
 - Open-file table: tracks open files
 - File pointer: pointer to last read/write location, per process that has the file open
 - File-open count: counter of number of times a file is open to allow removal of data from open-file table when last processes closes it
 - Disk location of the file: cache of data access information
 - Access rights: per-process access mode information

FILE LOCKING

- Provided by some operating systems and file systems
 - Similar to reader-writer locks
 - Shared lock similar to reader lock several processes can acquire concurrently
 - Exclusive lock similar to writer lock
- Mediates access to a file
- Mandatory or advisory:
 - Mandatory access is denied depending on locks held and requested
 - Advisory processes can find status of locks and decide what to do

FILETYPES - NAME, EXTENSION

file type	usual extension	function		
executable	exe, com, bin or none	ready-to-run machine- language program		
object	obj, o	compiled, machine language, not linked		
source code	c, cc, java, pas, asm, a	source code in various languages		
batch	bat, sh	commands to the command interpreter		
text	txt, doc	textual data, documents		
word processor	wp, tex, rtf, doc	various word-processor formats		
library	lib, a, so, dll	libraries of routines for programmers		
print or view	ps, pdf, jpg	ASCII or binary file in a format for printing or viewing		
archive	arc, zip, tar	related files grouped into one file, sometimes com- pressed, for archiving or storage		
multimedia	mpeg, mov, rm, mp3, avi	binary file containing audio or A/V information		

FILE STRUCTURE

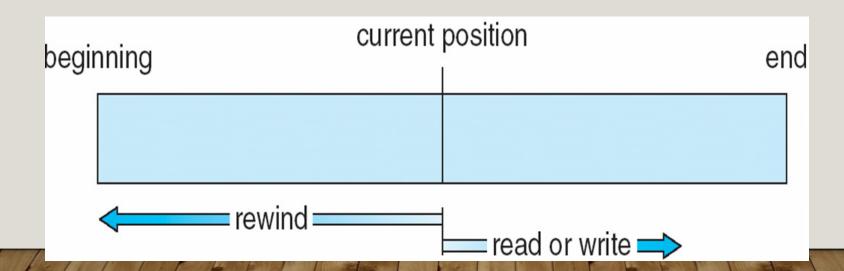
- Simple record structure
 - Lines
 - Fixed length
 - Variable length
- Complex Structures
 - Formatted document
 - Relocatable load file
- Can simulate last two with first method by inserting appropriate control characters
- Who decides:
 - Operating system
 - Program

ACCESS METHODS

- A file is fixed length logical records.
- The information in the file can be accessed in several ways:
- **▶** Sequential Access
- **▶ Direct Access**
- Other Access Methods

SEQUENTIAL ACCESS

- Information in the file is processed in order, one record after the other
- Operations
 - read next
 - write next
 - Reset
 - no read after last write (rewrite)



DIRECT ACCESS

- A file is made up of fixed-length logical records that allow programs to read and write records rapidly in no particular order.
- Based on a disk model of a file.
- Operations
 - read n
 - write n
 - position to n
 - read next
 - write next
 - rewrite n

n = relative block number

 Relative block numbers allow OS to decide where file should be placed

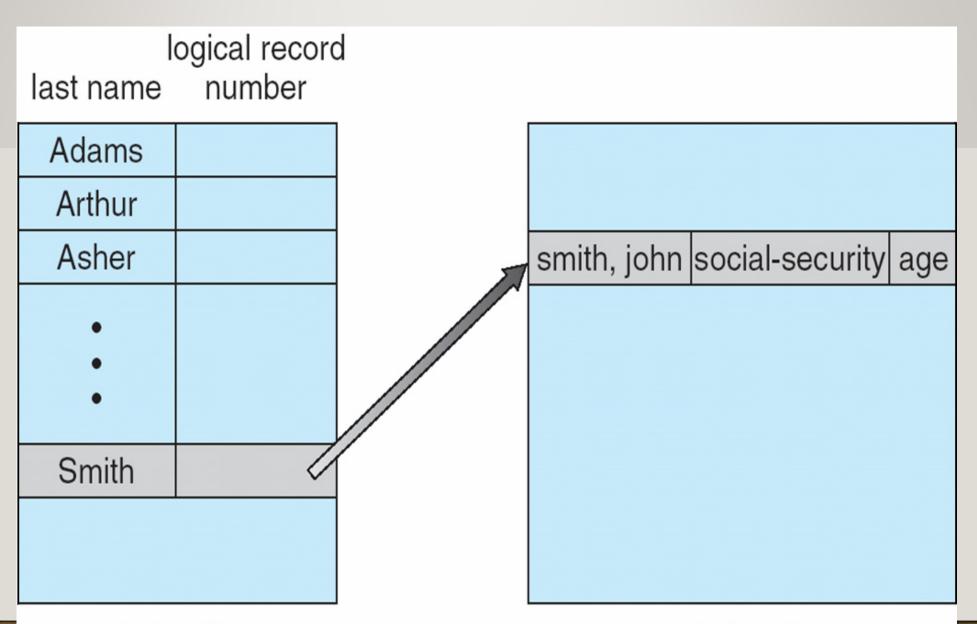
SIMULATION OF SEQUENTIAL ACCESS ON DIRECT-ACCESS FILE

sequential access	implementation for direct access		
reset	<i>cp</i> = 0;		
read next	read cp ; cp = cp + 1;		
write next	write cp ; cp = cp + 1;		

OTHER ACCESS METHODS

- Can be other access methods built on top of base methods
- General involve creation of an index for the file
- Keep index in memory for fast determination of location of data to be operated on (consider Universal Produce Code (UPC code) plus record of data about that item)
- If the index is too large, create an in-memory index, which is an index of a disk index
- IBM indexed sequential-access method (ISAM)
 - Small master index, points to disk blocks of secondary index
 - File kept sorted on a defined key
 - All done by the OS
- VMS operating system provides index and relative files as another example (see next slide)

EXAMPLE OF INDEX AND RELATIVE FILES

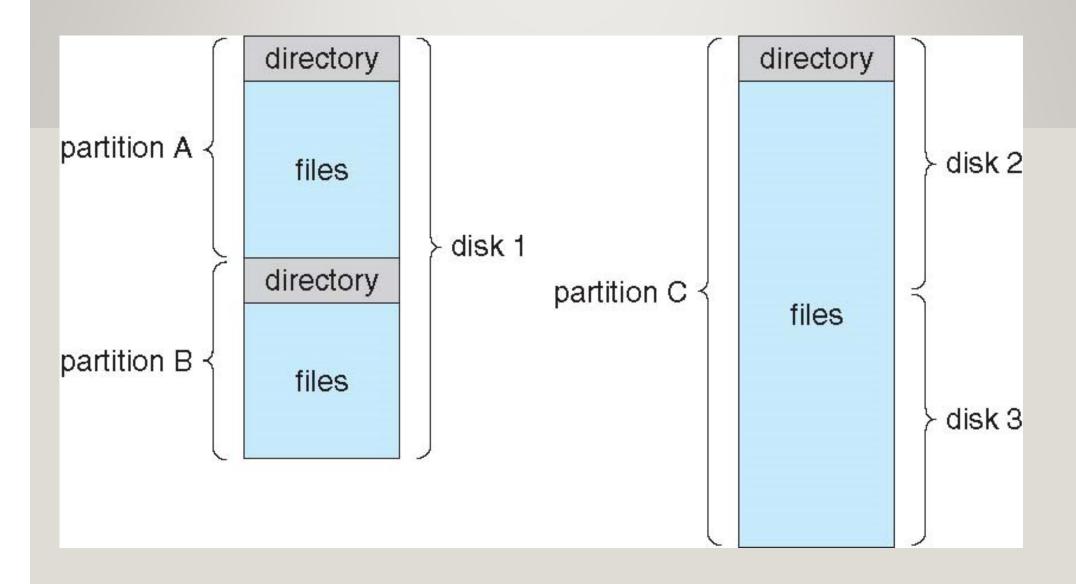


index file relative file

DISK STRUCTURE

- Disk can be subdivided into partitions
- Disks or partitions can be RAID protected against failure
- Disk or partition can be used raw without a file system, or formatted with a file system
- Partitions also known as minidisks, slices
- Entity containing file system is known as a volume
- Each volume containing a file system also tracks that file system's info in device directory or volume table of contents
- In addition to general-purpose file systems there are many special-purpose file systems, frequently all within the same operating system or computer

ATYPICAL FILE-SYSTEM ORGANIZATION

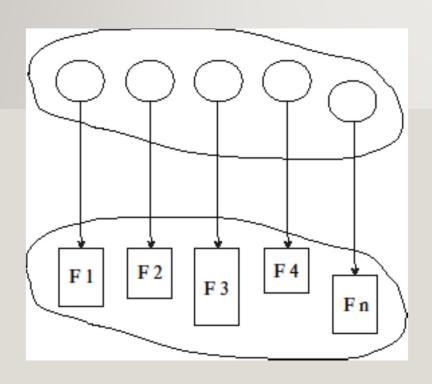


TYPES OF FILE SYSTEMS

- Typically, general-purpose file systems but systems may have some special- purpose file systems.
- Consider Solaris has
 - tmpfs memory-based volatile FS for fast, temporary I/O
 - objfs interface into kernel memory to get kernel symbols for debugging
 - ctfs contract file system for managing daemons
 - lofs loopback file system allows one FS to be accessed in place of another
 - procfs kernel interface to process structures
 - ufs, zfs general purpose file systems

DIRECTORY STRUCTURE

A collection of nodes containing information about all files



Operations performed on directory: search, create, delete, list, rename, traverse.

- Typically, a directory entry consists of the file's name and its unique identifier.
- Both the directory structure and the files reside on disk

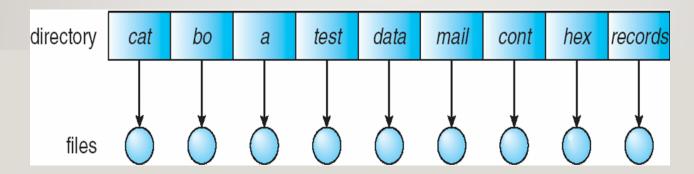
DIRECTORY ORGANIZATION

The directory is organized logically to obtain:

- Efficiency locating a file quickly
- Naming convenient to users
 - Two users can have same name for different files.
 - The same file can have several different names
- Grouping logical grouping of files by properties, (e.g., all Java programs, all games, ...)

SINGLE-LEVEL DIRECTORY

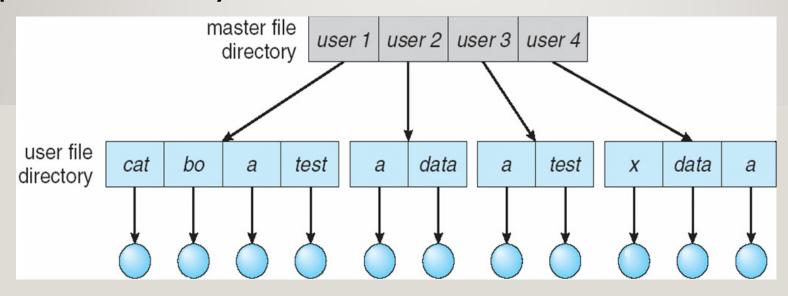
A single directory for all users



- Significant limitations, when the number of files increases or when the system has more than one user
- Naming problem
- Grouping problem

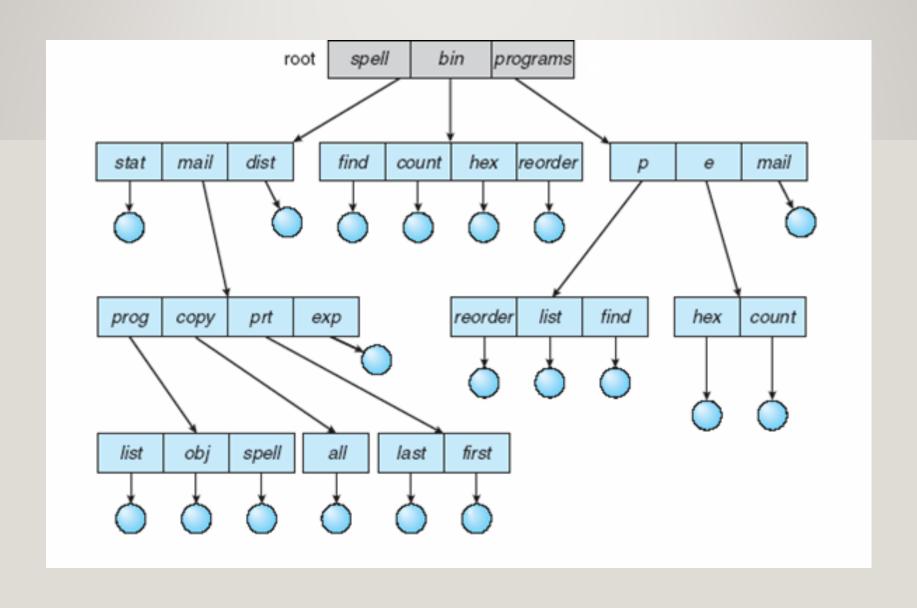
TWO-LEVEL DIRECTORY

Separate directory for each user



- Path name
- Can have the same file name for different user
- Efficient searching
- No grouping capability

TREE-STRUCTURED DIRECTORIES

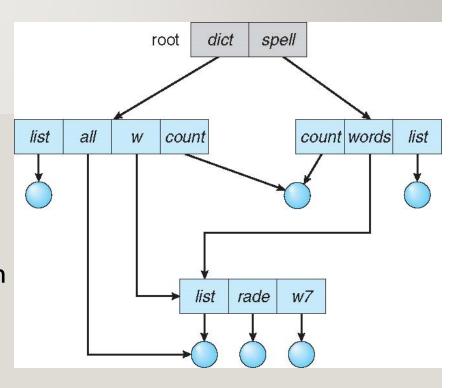


ACYCLIC-GRAPH DIRECTORIES

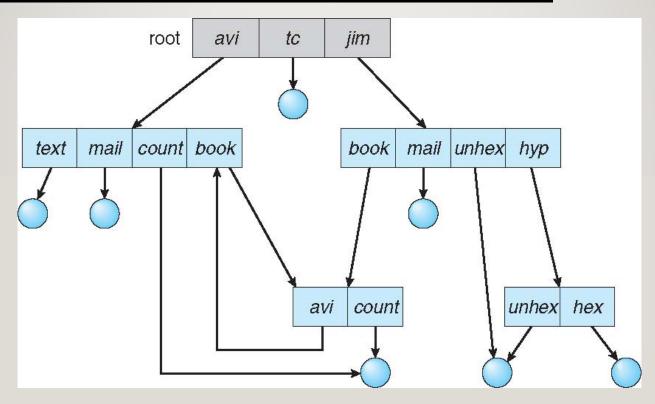
- Have shared subdirectories and files.
- Two different names (aliasing)
- If **dict** deletes $\mathbf{w}/\mathbf{list} \Rightarrow \text{dangling pointer}$

Solutions:

- Backpointers, so we can delete all pointers.
 - Variable size records a problem
- Backpointers using a daisy chain organization
- Entry-hold-count solution
- New directory entry type
 - Link another name (pointer) to an existing file
 - Resolve the link follow pointer to locate the file



GENERAL GRAPH DIRECTORY



How do we guarantee no cycles?

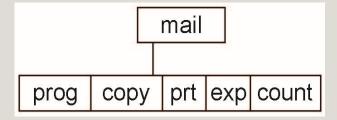
- > Allow only links to files not subdirectories
- > Garbage collection
- Every time a new link is added use a cycle detection algorithm to determine whether it is OK

CURRENT DIRECTORY

- Can designate one of the directories as the current (working) directory
 - cd /spell/mail/prog
 - type list
- Creating and deleting a file is done in current directory
- Example of creating a new file
 - If in current directory is /mail
 - The command

mkdir <dir-name>

Results in:



Deleting "mail" ⇒ deleting the entire subtree rooted by "mail"

PROTECTION

- File owner/creator should be able to control:
 - What can be done
 - By whom
- Types of access
 - Read
 - Write
 - Execute
 - Append
 - Delete
 - List

ACCESS LISTS AND GROUPS IN UNIX

- Mode of access: read, write, execute
- Three classes of users on Unix / Linux

a) owner access 7
$$\Rightarrow$$
 1 1 1 RWX
b) group access 6 \Rightarrow 1 1 0 RWX
c) public access 1 \Rightarrow 0 0 1

- Ask manager to create a group (unique name), say G, and add some users to the group.
- For a file (say game) or subdirectory, define an appropriate access.

Attach a group to a file

A SAMPLE UNIX DIRECTORY LISTING

-rw-rw-r	1 pbg	staff	31200	Sep 3 08:30	intro.ps
drwx	5 pbg	staff	512	Jul 8 09.33	private/
drwxrwxr-x	2 pbg	staff	512	Jul 8 09:35	doc/
drwxrwx	2 pbg	student	512	Aug 3 14:13	student-proj/
-rw-rr	1 pbg	staff	9423	Feb 24 2003	program.c
-rwxr-xr-x	1 pbg	staff	20471	Feb 24 2003	program
drwxxx	4 pbg	faculty	512	Jul 31 10:31	lib/
drwx	3 pbg	staff	1024	Aug 29 06:52	mail/
drwxrwxrwx	3 pbg	staff	512	Jul 8 09:35	test/

WINDOWS 7 ACCESS-CONTROL LIST MANAGEMENT

