Domain Name Systems

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Some of the slides borrowed from the book 'Computer Security: A Hands on Approach' by Wenliang Du

DNS Hierarchy

Lookup records for mapping from domain names to IP addresses



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Lookup records for mapping from domain names to IP addresses



Domain: Is a subtree, sharing its domain name with the name of the top most node in the subtree

DNS Hierarchy

Lookup records for mapping from domain names to IP addresses



SubDomains: Is a domain that branches off another.

Root Domain



Root Domain

List of Root Servers

HOSTNAME	IP ADDRESSES	MANAGER	
a.root-servers.net	198.41.0.4, 2001:503:ba3e::2:30	VeriSign, Inc.	10 in USA
b.root-servers.net	199.9.14.201, 2001:500:200::b	University of Southern California (ISI)	1 in Notherlands
c.root-servers.net	192.33.4.12, 2001:500:2::c	Cogent Communications	
d.root-servers.net	199.7.91.13, 2001:500:2d::d	University of Maryland	1 In Sweden
e.root-servers.net	192.203.230.10, 2001:500:a8::e	NASA (Ames Research Center)	1 in Japan
f.root-servers.net	192.5.5.241, 2001:500:2f::f	Internet Systems Consortium, Inc.	
g.root-servers.net	192.112.36.4, 2001:500:12::d0d	US Department of Defense (NIC)	Why only 13 root servers?
h.root-servers.net	198.97.190.53, 2001:500:1::53	US Army (Research Lab)	
i.root-servers.net	192.36.148.17, 2001:7fe::53	Netnod	567 mirrored root servers
j.root-servers.net	192.58.128.30, 2001:503:c27::2:30	VeriSign, Inc.	(9 mirrors in India – 2015)
k.root-servers.net	193.0.14.129, 2001:7fd::1	RIPE NCC	(3 I root servers: 2 I root serv
l.root-servers.net	199.7.83.42, 2001:500:9f::42	ICANN	1 L 1 K 1 E and D root sorvo
m.root-servers.net	202.12.27.33, 2001:dc3::35	WIDE Project	II, IK, IF, and Droot serve

https://internetdemocracy.in/wp-content/uploads/2016/03/Dr.-Anja-Kovacs-and-Rajat-Rai-Handa-India-at-the-Internets-Root.pdf 6



https://en.wikipedia.org/wiki/INRegistry

Top Level Domains

Top level domain

1547 as on July 2017



https://en.wikipedia.org/wiki/INRegistry

DNS Zone



Zone: Is a domain (or subdomain) that branches is served by a Name server.

A zone may be an entire domain with all its child domains, or a portion of a domain.

A zone can be the entire subtree starting at example.com

Or the company may decide to have several sub zones, for example one at usa.example.com

Authoritative Name Servers

[chester@aaha]	lya:~\$ nslookup
> set queryt	ype=soa
<pre>> iitm.ac.in</pre>	
Server:	10.6.0.11
Address:	10.6.0.11#53
Non-authorit	ative answer:
iitm.ac.in	
orig	in = dns1.iitm.ac.in
mail	addr = root.dns1.iitm.ac.in
seri	al = 2019020801
refr	esh = 10800
retr	y = 3600
expi	re = 1814400
mini	mum = 86400
Authoritativ	e answers can be found from:
iitm.ac.in	nameserver = dns2.iitm.ac.in.
iitm.ac.in	nameserver = dns1.iitm.ac.in.
dns1.iitm.ac	in internet address = 10.24.4.10.
dns2.iitm.ac	in internet address = 10.24.4.11.

→ Start of authority

Each DNS Zone has at least one authoritative name server that publishes information about that zone.

They are called `authoritative' because they provide original and answers to DNS queries as opposed to obtaining answers from other DNS servers.

2 authorities. 1 primary and the other secondary



Local DNS Server and Iterative Query



The iterative process starts from the ROOT Server. If it doesn't know the IP address, it sends back the IP address of the nameservers of the next level server (.NET server) and then the last level server (example.net) which provides the answer.

http://www.iitm.ac.in — Entered in web-browser



Local System:

Lookup /etc/hosts file.

Can the /etc/hosts file resolve (have the IP address) for www.iitm.ac.in?

chester@aahalya	:~\$ cat /etc/hosts	
127.0.0.1	localhost	
10.6.15.91	aahalya.cse.iitm.ac.in	aahalya
10.6.15.92	aampal.cse.iitm.ac.in	aampal
10.6.15.93	aarunya.cse.iitm.ac.in	aarunya
10.6.15.94	aarika.cse.iitm.ac.in	aarika
10.6.15.95	baarika.cse.iitm.ac.in	baarika
10.6.15.96	basavi.cse.iitm.ac.in	basavi
10.6.15.97	cabitha.cse.iitm.ac.in	cabitha
chester@aahalya	:~\$	

http://www.iitm.ac.in — Entered in web-browser

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Local DNS Server: Lookup the local DNS server (server present in the LAN).

How to identify the IP address of the Local DNS server? (/etc/resolv.conf) This needs to be configured or, can be found, if the system is configured for DHCP, then this file is automatically modified.

```
chester@aahalya:~$ cat /etc/resolv.conf
search cse.iitm.ac.in
nameserver 10.6.0.11
nameserver 10.6.0.12
```

If the local DNS server can resolve the address; then we are done. Else, the resolver would be activated. The resolver would need to query another DNS server, higher up in the hierarchy.

http://www.iitm.ac.in.

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Resolver in Local DNS will query the Root Name Server

→(from resolver) What is the IP address of www.iitm.ac.in

← (from root server)I don't know the answer, you can ask any of these authorities

chester@aahalya:~\$ dig	@a.root-	servers.	net iitm	.ac.in
;; QUESTION SECTION:				
;iitm.ac.in.		IN	Α	
:: AUTHORITY SECTION:				
in.	172800	IN	NS	a0.in.afilias-nst.info.
in.	172800	TN	NS	a1.in.afilias-nst.in.
in.	172800	TN	NS	a2.in.afilias-nst.info.
in.	172800	IN	NS	b0.in.afilias-nst.org.
in.	172800	TN	NS	b1.in.afilias-nst.in.
in.	172800	TN	NS	b2.in.afilias-nst.org.
in.	172800	IN	NS	c0.in.afilias-nst.info.
;; ADDITIONAL SECTION:				
a0.in.afilias-nst.info.	172800	IN	Α	199.7.87.1
a1.in.afilias-nst.in.	172800	IN	Α	115.249.164.142
a2.in.afilias-nst.info.	172800	IN	Α	199.249.117.1
o0.in.afilias-nst.org.	172800	IN	Α	199.253.56.1
o1.in.afilias-nst.in.	172800	IN	Α	180.179.215.70
o2.in.afilias-nst.org.	172800	IN	Α	199.249.125.1
c0.in.afilias-nst.info.	172800	IN	Α	199.253.57.1
a0.in.afilias-nst.info.	172800	IN	AAAA	2001:500:29::1
a1.in.afilias-nst.in.	172800	IN	AAAA	2001:4528:fff:13::142
a2.in.afilias-nst.info.	172800	IN	AAAA	2001:500:45::1
o0.in.afilias-nst.org.	172800	IN	AAAA	2001:500:50::1
o1.in.afilias-nst.in.	172800	IN	AAAA	2401:8800:411:8::70
o2.in.afilias-nst.org.	172800	IN	AAAA	2001:500:4d::1
c0.in.afilias-nst.info.	172800	IN	AAAA	2001:500:51::1

Directly send the query to this server.

[chester@aahalya:~\$ dig @a0.in.afilias-nst.info iitm.ac.in

http://www.iitm.ac.in.

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Resolver in Local DNS will query the TLD

→(from resolver) What is the IP address of <u>www.iitm.ac.in</u>

← (return)I don't know the answer, you can ask any of these authorities

;; QUESTION SECTION: ;iitm.ac.in.		IN	A	
;; AUTHORITY SECTION:				
iitm.ac.in.	86400	IN	NS	dns3.iitm.ac.in.
iitm.ac.in.	86400	IN	NS	dns2.iitm.ac.in.
iitm.ac.in.	86400	IN	NS	dns1.iitm.ac.in.
;; ADDITIONAL SECTION:				
dns1.iitm.ac.in.	86400	IN	Α	203.199.213.2
dns2.iitm.ac.in.	86400	IN	Α	14.130.160.3
dns3.iitm.ac.in.	86400	IN	Α	14.139.160.2

http://www.iitm.ac.in.

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Resolver in Local DNS will query the next level NS

→What is the IP address of <u>www.iitm.ac.in</u>
 ← The SOA is dns1.iitm.ac.in

[chester@aahalya:~\$ dig (@dns1.iitm.ac	.in iitm.	ac.in		
; <<>> DiG 9.7.3 <<>> Qa ; (1 server found) ;; global options: +cmd ;; Got answer: ;; ->>HEADER<<- opcode: ;; flags: qr aa rd ra; Q	dns1.iitm.ac.: QUERY, statu: QUERY: 1, ANSN	in iitm.a s: NOERRO WER: 0, A	c.in R, id: 15832 UTHORITY: 1, /	ADDITIC	DNAL: 0
;; QUESTION SECTION: ;iitm.ac.in.	IN	A			
;; AUTHORITY SECTION: iitm.ac.in. ac.in. 2019020801 10800	86400 IN 3600 1814400	SOA 86400	dns1.iitm.a	ac.in.	root.dns

DNS Cache

- The local DNS server will cache all responses from other DNS servers (to reduce queries)
- TTL available with all responses, which determines when the entry would be removed from cache



DNS Cache

- Due to caching
 - Most resolver queries do not need a query to the root server
 - 2% of all queries to the root-servers are legitimate
 - 75% were due to incorrect or non-existent caching
 - 12.5% to unknown TLDs
 - 7% were for lookups to IP addresses, as if, they were domain names

Set Up DNS Zones on Local DNS Server

- > Utility in Linux: bind9
- Create zones: Create two zone entries in the DNS server by adding them to /etc/bind/named.conf.

```
zone "example.net" {
    type master;
    file "/etc/bind/example.net.db";
    };
zone "0.168.192.in-addr.arpa" {
    type master;
    file "/etc/bind/192.168.0.db";
    };
For reverse lookup
(IP → hostname).
```

Zone File for Forward Lookup

/etc/bind/example.net.db (The file name is specified in named.conf)

: their own TTL	s the
a IN SOA ng ayampla nat admin ayampla nat (a) Represent	s the
e in sox istexample.net. admintexample.net. (e.nepiesent	
1 ; Serial origin specific	d in ا
8H ; Refresh Oligin Specific	
2H ; Retry named.con	f (string
4W ; Expire	
1D) ; Minimum alter 2016 j	
[example.r	net]
<pre>@ IN NS ns.example.net. ;Address of nameserver</pre>	-
@ IN MX 10 mail.example.net. ;Primary Mail Exchanger	
www IN A 192.168.0.101 ;Address of www.example.net	
mail IN A 192.168.0.102 ;Address of mail.example.net	
ns IN A 192.168.0.10 ;Address of ns.example.net	
*.example.net. IN A 192.168.0.100 ;Address for other URL in	
; the example.net domain	

Zone File for Reverse Lookup

/etc/bind/192.168.0.db: (The file name is specified in named.conf)

\$TTL 3D				
Q	IN	SOA 1 8H 2H 4W 1D)	ns.example.net. admin.example.net.	
0	IN	NS	ns.example.net.	
101	IN	PTR	www.example.net.	
102	IN	PTR	mail.example.net.	
10	IN	PTR	ns.example.net.	

Testing the setup

\$ dig www.example.net <<>> DiG 9.5.0b2 <<>> www.example.net ;; global options: printcmd ;; Got answer: ;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 27136 ;; flags: gr aa rd ra; QUERY: 1, ANSWER: 1, AUTHORITY: 1, ADDITIONAL: 1 ;; QUESTION SECTION: ;www.example.net. ΤN А ;; ANSWER SECTION: www.example.net. A 192.168.0.101 259200 IN ;; AUTHORITY SECTION: example.net. NS ns.example.net. 259200 IN ;; ADDITIONAL SECTION: ns.example.net. 259200 IN A 192.168.0.10

Need to ensure that Resolv.conf is pointing to the recently setup DNS server

DNS Queries

Message formats are defined in RFC 1035.

The good news is that each message has the same generic format with 5 sections. This is the last good news.

Section	Meaning/Use
Section 1	Message Header
Section 2	The DNS question being asked (aka Question Section)
Section 3	The Resource Record(s) which answer the question (aka Answer Section)
Section 4	The Resource Record(s) which point to the domain authority (aka Authority Section)
Section 5	The Resource Record(s) which may hold additional information (aka Additional Section)

DNS Query Format

Message Header

0		1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
Message ID																
Q	R	O	PC	PCODE AA TC RD RA res1 res2 res3 RCODE												
				QDCOUNT (No. of items in Question Section)												
				A	NC	οµ	NT (I	No. o	f iter	ns in <i>i</i>	Answe	er Sec	tion)		
				NS	CC	οŲΝ	IT (N	lo. of	item	ns in A	uthor	ity Se	ctior	ר)		
			/	٩R	СС	NUN	T (N	o. of	item	s in A	dditio	nal Se	ectio	n)		
	Authoritative Answer															
	 • 0: query, 1: inverse query; 2: status 															
С	2R	=() c	านเ	er	y; (QR=	1 re	espo	onse						

Sent in the query and reflected back by the response

http://www.zytrax.com/books/dns/ch15/

DNS Question Section

Question S	Section	/ assume // the he: 03 77 77 // printal ! w w	an A query with a name of www.mydomain.com x representation is 77 08 6D 79 64 6F 6D 61 69 6E 03 63 6F 6D 00 ble w ! m y d o m a in ! c o m !
Field Name	Meaning/Use		
QNAME	The domain name being queried		
QTYPE	The resource records being requested	Value	Meaning/Lise
QCLASS	The Resource Record(s) class being requested, for instance, inter	value	Requests the A record for the domain name
		x'0002 (2)	Requests the NS record(s) for the domain name
		x'0005 (5)	Requests the CNAME record(s) for the domain name
		x'0006 (6)	Requests the SOA record(s) for the domain name
		x'000B (11)	Requests the WKS record(s) for the domain name
		x'000C (12)	Requests the PTR record(s) for the domain name
		x'000F (15)	Requests the MX record(s) for the domain name
		x'0021 (33)	Requests the SRV record(s) for the domain name
		x'001C (28)	Requests the AAAA record(s) for the domain name
//www.zytra	ax.com/books/dns/ch15/	x'00FF (255)	Requests ANY resource record (typically wants SOA, MX, NS and MX)

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DNS Question Section

Answer Section

Field Name	Meaning/Use
NAME	The name being returned e.g. www or ns1.example.net If the name is in the same domain as the question then typically only the host part (label) is returned, if not then a FQDN is returned.
TYPE	The RR type, for example, SOA or AAAA
CLASS	The RR class, for instance, Internet, Chaos etc.
TTL	The TTL in seconds of the RR, say, 2800
RLENGTH	The length of RR specific data in octets, for example, 27
RDATA	The RR specific data (see Binary RR Formats below) whose length is defined by RDLENGTH, for instance, 192.168.254.2

DNS Question Section

Authority Section

Field Name	Meaning/Use
NAME	The name being returned e.g. www or ns1.example.net If the name is in the same domain as the question then typically only the host part (label) is returned, if not then a FQDN is returned.
TYPE	The RR type, for example, SOA or AAAA
CLASS	The RR class, for instance, Internet, Chaos etc.
TTL	The TTL in seconds of the RR, say, 2800
RLENGTH	The length of RR specific data in octets, for example, 27
RDATA	The RR specific data (see Binary RR Formats below) whose length is defined by RDLENGTH, for instance, 192.168.254.2

This section mentions the servers that are the ultimate authority for answering DNS queries.

Answers, may be obtained from the cache of other DNS servers.

Can be used to check with the authoritative response.









Local DNS Cache Poisoning Attack

Goal: Forge DNS replies after seeing a query from Local DNS Server **Technique:** Sniffing and Spoofing

Local DNS Cache Poisoning Attack



\$ dig www.example.net

Result

;; QUESTION SECTION: ;www.example.net.		IN	А	
;; ANSWER SECTION: www.example.net.	19000	IN	A	10.20.30.40
;; AUTHORITY SECTION:	19000	IN	NS	ns.example.net.
;; ADDITIONAL SECTION: ns.example.net.	19000	IN	A	10.20.30.50

Attack

Inspect the Cache

;				
\$DATE 201512141528	354			
; authanswer				
	18997	IN	NS	ns.example.net.
; authauthority				
ns.example.net.	18997	NS		ns.example.net.
; additional				
	18997	А		10.20.30.50
; authanswer				
www.example.net.	18997	А		10.20.30.40
;				

- Run"sudo rndc dumpdb -cache" and check the contents of "/var/cache/bind/dump.db".
- Clean the cache using "sudo rndc flush" before doing the attack.

Targeting the Authority section



Any DNS query sent to the local DNS server will be (if needed) directed to the attacker's ns.attacker.net



Cache Poisoning Without Sniffing

Section	Meaning/Use													
Section 1	Message Header													
Section 2	The DNS question being asked (aka Question Section	0	1 2	3 4	5	6	7	8	9	10	11	12 1	13 1	4 15
Section 3	The Resource Record(s) which answer the question Message ID							ge ID						
Section 4	The Resource Record(s) which point to the domain a	QR	OPC	ODE		ТС	RD	RA	res1	res2	res3	RCC	DE	
Section 5	The Resource Record(s) which may hold additional i QDCOUNT (No. of items in C							s in C	Questi	on Se	ction)			
	·			ANC	OU	NT (N	lo. of	fiten	ns in <i>I</i>	Answe	er Sec	ction)		
		NSCOUNT (No. of items in Authority Section)												
		ARCOUNT (No. of items in Additional Section)												

Two difficulties in creating a valid spoof:

- 1. Need to guess the local DNS server's source port (2^16 possibilities)
- 2. The response should have the same Message ID as the DNS query in step (2).

(Brute force attack $2^{32} \rightarrow$ at 1000 spoofed queries / second, it will take 50 days to try all 2^32 possibilities

Further Difficulties: the Local DNS server's cache

If the real response (3) arrives and it is cached (4). Then subsequent queries will read off the cache (5 \rightarrow 6 \rightarrow 7) and no query is made from the Local DNS.

Thus, to make another try, the attacker should wait till the cache is flushed.



Flaw in the Protocol

- When looking up sibling names like 1.google.com, and 2.google.com.
 - Attackers can do this and say they're the official server for www.google.com, telling the local DNS server what www. needs to be, and the local DNS will believe the attacker.

Kaminsky Attack

How to keep trying spoofed DNS responses (2^32 times) without worrying about the cache effect?

Kaminsky's Idea:

- Ask a different question every time, so caching the answer does not matter, and the local DNS server will send out a new query each time.
- Provide forged answer in the Authority section

Kaminsky Attack



The Kaminsky Attack: A Sample Response



Spoofing Replies: IP and UDP headers



Spoofing Replies: DNS Header and Payload

Question Record

Name	Record Type	Class			
twysw.example.com	"A" Record 0x0001	Internet 0x0001			

Answer Record

Name	Record Type	Class	Time to Live	Data Length	Data: IP Address
twysw.example.com	"A" Record 0x0001	Internet 0x0001	0x00002000 (seconds)	0x0004	1.2.3.4

Authority Record

Name	Record Type	Class	Time to Live	Data Length	Data:	Name Server
example.com	"NS" Record 0x0002	Internet 0x0001	0x00002000 (seconds)	0x0013	ns.at	tacker32.net
			2 n s 10 a	Represent (Tota	tation in the packet al: 0x13 bytes) c k e r 3 2	3 c o m 0

Digital Signatures



Signing Function

 $y = sig_a(x)$

Input : Message (x) and Alice's private key **Output:** Digital Signature of Message

Verifying Function

ver_b(x, y)

Input : digital signature, message Output : true or false true if signature valid false otherwise

Protection Against DNS Cache Poisoning Attacks

DNSSEC

- DNSSEC is a set of extension to DNS, aiming to provide authentication and integrity checking on DNS data.
- With DNSSEC, all answers from DNSSEC protected zones are digitally signed.
- By checking the digital signatures, a DNS resolver is able to check if the information is authentic or not.
- DNS cache poisoning will be defeated by this mechanism as any fake data will be detected because they will fail the signature checking.

Protection Using DNSSEC



Protection Using TLS/SSL

Transport Layer Security (TLS/SSL) protocol provides a solution against the cache poisoning attacks.

- After getting the IP address for a domain name (<u>www.example.net</u>) using DNS protocol, a computer will ask the owner (server) of the IP address to prove that it is indeed www.example.net.
- The server has to present a public-key certificate signed by a trusted entity and demonstrates that it knows the corresponding private key associated with www.example.net (i.e., it is the owner of the certificate).
- HTTPS is built on top of TLS/SSL. It defeats DNS cache poisoning attacks.

DNSSEC versus TLS/SSL

- Both DNSSEC and TLS/SSL are based on the public key technology, but their chains of trust are different.
- DNSSEC provides chain of trust using DNS zone hierarchy, so nameservers in the parent zones vouch for those in the child zones.
- TLS/SSL relies on Public Key Infrastructure which contains Certificate Authorities vouching for other computers.

Denial of Service Attacks on Root Servers

Attacks on the Root and TLD Servers :

<u>Root nameservers:</u> If the attackers can bring down the servers of the root zone, they can bring down the entire Internet. However, attack root servers is difficult:

- The root nameservers are highly distributed. There are 13 (A,B....M) root nameservers (server farm) consisting of a large number of redundant computers to provide reliable services.
- As the nameservers for the TLDs are usually cached in the local DNS servers, the root servers need not be queried till the cache expires (48 hrs). Attacks on the root servers must last long to see a significant effect.

Denial of Service Attacks on TLD Servers

Nameservers for the TLDs are easier to attack. TLDs such as gov, com, net etc have quite resilient infrastructure against DOS attacks. But certain obscure TLDs like country-code TLDs do not have sufficient infrastructure. Due to this, the attackers can bring down the Internet of a targeted country.

Attacks on Nameservers of a Particular Domain

Dyn network : In 2016, multiple DDoS attacks were launched against a major DNS service provider for companies like CNN, BBC, HBO, PayPal etc. The attacks are believed to have been launched through botnet consisting of different IoT devices like IP cameras, baby monitors etc. It caused major Internet services unavailable .



Gizmodo @ @Gizmodo · Oct 22 Yesterday's brutal DDoS attack is the beginning of a bleak future gizmo.do/POr2Sne



Mirai Botnet Malware

- Number of IoT devices increasing at a rapid rate
- These devices are characterized by
 - Low profile
 - Less user interactions

Low hanging fruit for hackers

- Security often compromised (for better performance / smaller profile)
- Not always up-to-date with security patches
- Malware with IoT devices as targets
 - Bashlight and Mirai are the most popular
 - PNScan, targets x86 platforms.
 - Try to determine router login baed on a special dictionary
 - Connect using ssh connection using predefined user credentials

Mirai BotNet Malware

280 Gbps max flooding 50,000 unique lps 164 countries







Mirai BotNet Malware





The Mirai Botnet and the IoT Zombie Armies, 2017



Mirai BotNet Malware



Newly formed bot establishes connection with C&C. Periodic heartbeats between the two.

Other activities in the bot:

* memory scraping to identify other malware present in the bot. If found, kill the process. (wants to be the own the device)

- Deletes itself from the persistent storage. Will only be available in RAM (fileless malware)
- Monitors the watchdog timer to defend against system hangs and reboots.
- On command, can start a variety of floods: SYN, ACK, UDP, GRE, DNS, STOMP, ETH. Application layer flooding. 25,000 SYN packets per second.

Mirai BotNet Malware



Other Mirai Variants

- US university (Feb 2017)
- Windows based strain
- Mirai strain used for bitcoin mining

Detection

- Patterns of communications
 - Huge communication on ports 23, 2323, 22 for authorization purposes
 - Frequent exchange of traffic with infrastructure
 - Surge of egress traffic throughout the course of the attack



Mitigation

Short-term mitigations

- Blocking TCP ports used for probing and bruteforcing the device
- Filtering egress and ingress packets by TCP rules

Better design strategies

- Least privilege implementations
- capability based systems
- Microkernel /Unikernel based approaches (Linux is too large for embedded applications)