

Categories and Objects

- Organizing objects into Categories is fundamental to Knowledge Representation.
- Interactions take place at object level but the reasoning takes place at the category level.
 - A shopper goes to a shop to buy a basketball
(as supposed to buying the 3rd basketball on the 4th shelf)
- Agent infers the object using percepts then uses the category information to make prediction about the object
 - If Agent detects yellow peel, yellow soft pulp and a hard shell inside then Agent infers that the object is Mango
 - Using its category information, it infers that it can used to make milkshake

Categories and Objects

- There are two ways to represent categories and objects in First Order Logic:
 - Make Categories as unary predicates
 - `BasketBall(B1)`
 - Make the category also as an object (Reification)
 - `Member(B1, BasketBalls)` also written as `B1 ∈ BasketBalls`
 - `SubClass(BasketBalls, Balls)` also written as `BasketBalls ⊆ Balls`
- Categories are organized through inheritances
 - Subclasses organize categories into taxonomic hierarchy
 - Examples:
 - Library Dewey Decimal system
 - Biological Taxonomy : Kingdom, Phylum, Class, Order, Family, Genus, Species

Categories in First order logic

- Many natural properties of Categories can be stated in First Order Logic
 - An object is a member of a category
 - $B1 \in \text{BasketBalls}$
 - A category is a subclass of another category
 - $\text{BasketBalls} \subseteq \text{Balls}$
 - Members of a class can be recognized by some properties
 - $\text{Orange}(x) \wedge \text{Spherical}(x) \wedge \text{Diameter} = 9.5'' \wedge x \in \text{Balls} \Rightarrow x \in \text{BasketBalls}$
 - All members of a category have some common property
 - $\text{Dogs} \in \text{DomesticatedSpecies}$
 - Here DomesticatedSpecies is a category of categories

Categories in First order logic

- Sometimes we also want to say two categories are disjoint / exhaustive / partition:
 - $\text{Disjoint}(\{\text{Animals}, \text{Vegetables}, \text{Basketballs}\})$
 - $\text{Disjoint}(s) : \forall c1, c2 \quad c1 \in s \wedge c2 \in s \wedge c1 \neq c2 \Rightarrow \text{Intersection}(c1, c2) = \{ \}$
 - $\text{ExhaustiveDecomposition}(\{\text{CSE}, \text{AE}, \text{EE}, \dots\}, \text{BTechStudents})$
 - $\text{ExhaustiveDecomposition}(s, c) : \forall i \quad i \in c \Leftrightarrow \exists c2 \quad c2 \in s \wedge i \in c2$
 - $\text{Partition}(\{\text{BTech}, \text{DualDegree}, \text{MTech}, \text{MS}, \text{PhD}\}, \text{Students})$
 - $\text{Partition}(s, c) : \text{Disjoint}(s) \wedge \text{ExhaustiveDecomposition}(s, c)$
- Some categories can also have necessary and sufficient conditions for membership
 - $x \in \text{Bachelors} \Leftrightarrow \text{Unmarried}(x) \wedge x \in \text{Adults} \wedge x \in \text{Males}$

Representing Physical Composition

- Nose is **part of** Face
 - Chennai is **part of** Tamil Nadu
 - Tamil Nadu is **part of** India
 - India is **part of** Asia
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- Objects can be grouped into **Partof** hierarchies
 - **Partof(x,y)** is reflexive and transitive
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- Category of composite objects can be characterized using **PartOf**
 - **Archipelago is a collection of at least 2 islands**
 - $x \in \text{Archipelago} \Leftrightarrow \exists I1\ I2\ I1 \neq I2 \wedge \text{PartOf}(I1, x) \wedge \text{PartOf}(I2, x) \wedge I1 \in \text{island} \wedge I2 \in \text{island}$
-
- $$\wedge \forall I\ \text{PartOf}(I, x) \Rightarrow I \in \text{island}$$

Representing Measurements

- Objects have **height / mass / cost ...**
- Same length has different measurements (inches / cm / ..)
 - **Each unit can be a function**
 - `length(l1) = Inches(1.5)` `length(l1) = centimeters(3.8)`
- Examples of simple measurements:
 - `Diameter(BasketBall1) = Inches(9.5)`
 - `MRP(BasketBall1) = INR(1200)`
 - `d ∈ days ⇒ (duration(d) = hours(24))`
- **Inches(0)** and **Centimeters(0)** refer to the same object, but different from `seconds(0)`

Representing Measurements

- Not clear how to represent measurement that do not have agreed scale of values
 - Tasty / difficult / scared
- Does not make sense to impose numerical scale
 - Does it mean we cannot do any inference about it?
 - We are generally interested in the relative ordering, not the absolute value
 - As long as we can define $>$ we can make inferences
 - Fried food is tastier than salad
 - $\forall f1, f2 \quad f1 \in \text{FriedFood} \wedge f2 \in \text{salad} \Rightarrow \text{Tastyness}(f1) > \text{Tastyness}(f2)$
 - Tasty things are unhealthy
 - $\forall f1, f2 \quad \text{Tastyness}(f1) > \text{Tastyness}(f2) \Rightarrow \text{HealthScore}(f1) < \text{healthscore}(f2)$
- Such inferences in AI is called qualitative physics
 - Reasoning about physical systems without looking at detailed equations and numerical simulations

Representing Objects

- Real world consists of primitive objects
 - Composite objects are built from primitive objects
 - Part of is useful in many ontologies
 - There are objects that cannot be broken into primitive objects
 - Such objects are called Stuff
 - If I have a banana and butter in front of me:
 - I can say I see 1 banana (Banana is a thing)
 - Can I say I see 1 butter? (Butter is stuff)
 - There are ontologies to deal with this distinction
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Representing Objects

- For **stuff**, we can have a category
 - $\forall p, b \quad b \in \text{Butter} \wedge \text{PartOf}(p, b) \Rightarrow p \in \text{Butter}$
 - $\forall b \quad b \in \text{Butter} \Rightarrow \text{MeltingPoint}(b) = \text{centigrade}(30)$
- **PoundofButter** is **not stuff**, it is a **thing**
- **SalterButter** is **stuff** and is a subcategory of **Butter**
- **Intrinsic properties** like density, melting point etc **can be attributed to stuff**
- **Extrinsic properties** like weight, length etc **can be attributed to things**

Events

- We had fluents in the Wumpus world like HaveArrow^t
 - Here **time is discrete**
 - How to model **continuous time**?
 - Like **Water filling a bucket**:
 - Initial state : Empty bucket Final state : Full bucket
 - Transitions?
 - How to model **simultaneous events**?
 - **The boy was brushing his teeth while waiting for the bucket to fill**
 - To handle this, we have **Event calculus**
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Event Calculus

- Objects in event calculus are Events, Fluents and Time points
 - $At(Alice, Chennai)$ is a fluent that states Alice is in Chennai
- Alice flew from Chennai to Mumbai
 - Option 1: Make this a predicate (of appropriate arity)
 - Option 2:
 - $E1 \in Flyings \wedge Flyer(E1, Alice) \wedge Origin(E1, Chennai) \wedge Destination(E1, Mumbai)$
 - Since events are also objects, we can add more properties about events:
 - The flight was delayed by 1 hour: $delayed(E1) = hour(1)$
- Continuous time:
 - $T(At(Alice, Chennai), t1, t2)$ T is special : It takes a predicate as input
 - $Happens(E1, t1, t2)$
- What all such predicates do we need to model time/events?
 - Starts / terminates / ...

Event Calculus

- Predicates in Event Calculus (1999, Shanahan)

- $T(f, t1, t2)$ Fluent f is true for all times between $t1$ and $t2$
- $Happens(e, t1, t2)$ Event e starts at time $t1$ and ends at time $t2$
- $Initiates(e, f, t)$ Event e causes Fluent f to become true at time t
- $Terminates(e, f, t)$ Event e causes Fluent f to become false at time t
- $Initiated(f, t1, t2)$ Fluent f become(s) true at some point between $t1$ and $t2$
- $Terminated(f, t1, t2)$ Fluent f become(s) false at some point between $t1$ and $t2$
- $t1 < t2$ Time point $t1$ occurs before $t2$

- Example:

- $E1 \in \text{Flyings} \wedge \text{Flyer}(E1, \text{Alice}) \wedge \text{Origin}(E1, \text{Chennai}) \wedge \text{Destination}(E1, \text{Mumbai}) \wedge \text{Happens}(E, t1, t2) \Rightarrow$
 $\text{Terminates}(E, \text{At}(\text{Alice}, \text{Chennai}), t1) \wedge \text{Initiates}(E, \text{At}(\text{Alice}, \text{Mumbai}), t2)$

Event Calculus

- There is a distinguished **start event**
 - Describes **initial state** and **fluents that are true/false at the beginning**
 - We can then describe what fluent are true/false at what points of time
- Example:
 - Suppose an event **e** happens between **t1** and **t3**. At time **t2** somewhere between **t1** and **t3**, the event **e** changes the value of the fluent **f** by initiating it. Then at time **t4** in the future (after **t3**), if no other event has changed the value of **f** then the fluent **f** remains True between **t2** and **t4**.
 - $\text{Happens}(e, t1, t3) \wedge \text{Initiates}(e, f, t2) \wedge \neg \text{Terminated}(f, t2, t4) \wedge t1 \leq t2 \leq t3 \leq t4 \Rightarrow T(f, t2, t4)$
- Event calculus can be extended to include:
 - Simultaneous events (two children playing see-saw)
 - Exogenous event (wind moving an object)
 - Continuous event (rise of tides)
 - Non-deterministic events (pressing a button on coffee vending machine)

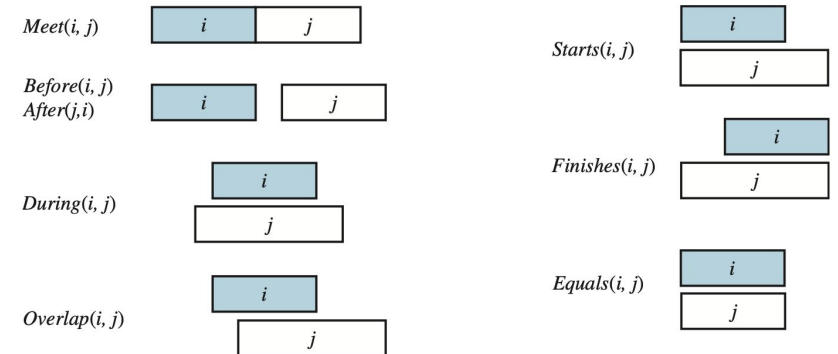
Time

- Event calculus allows us to talk about time points and time intervals
 - Time points are special time intervals with 0 duration
 - $\text{Partition}(\{\text{Moments}, \text{ExtendedIntervals}\}, \text{Intervals})$
 $i \in \text{Moments} \Leftrightarrow \text{Duration}(i) = \text{Seconds}(0)$
- How to indicate absolute time?
 - Start from an arbitrary absolute time point as 0
 - January 1, 1900 (GMT) has time 0
 - **Begin** and **End** functions pick the earliest and latest moments in an interval
 - **Time** function gives a point on the time scale
- Examples:
 - $\text{Time}(\text{Begin}(\text{AD1900})) = \text{Seconds}(0)$
 - $\text{Time}(\text{Begin}(\text{AD2001})) = \text{Seconds}(3187324800)$
 - $\text{Time}(\text{End}(\text{AD2001})) = \text{Seconds}(3218860800)$
 - $\text{Interval}(i) \Rightarrow \text{Duration}(i) = (\text{Time}(\text{End}(i)) - \text{Time}(\text{Begin}(i)))$
 - $\text{Duration}(\text{AD2001}) = \text{Seconds}(31536000)$
 - $\text{Time}(\text{Begin}(\text{AD2001})) = \text{Date}(0, 0, 0, 1, \text{Jan}, 2001)$
 - $\text{Date}(0, 20, 21, 24, 1, 1995) = \text{Seconds}(3000000000)$

Time Intervals

- Two events can interact in the following 7 ways (Allen, 1983)

- Meet(i, j) $\Leftrightarrow \text{End}(i) = \text{Begin}(j)$
- Before(i, j) $\Leftrightarrow \text{End}(i) < \text{Begin}(j)$
- After(j, i) $\Leftrightarrow \text{Before}(i, j)$
- During(i, j) $\Leftrightarrow \text{Begin}(j) < \text{Begin}(i) < \text{End}(i) < \text{End}(j)$
- Overlap(i, j) $\Leftrightarrow \text{Begin}(i) < \text{Begin}(j) < \text{End}(i) < \text{End}(j)$
 - Not symmetric : i should begin before j
- Starts(i, j) $\Leftrightarrow \text{Begin}(i) = \text{Begin}(j)$
- Finishes(i, j) $\Leftrightarrow \text{End}(i) = \text{End}(j)$
- Equals(i, j) $\Leftrightarrow \text{Begin}(i) = \text{Begin}(j) \wedge \text{End}(i) = \text{End}(j)$



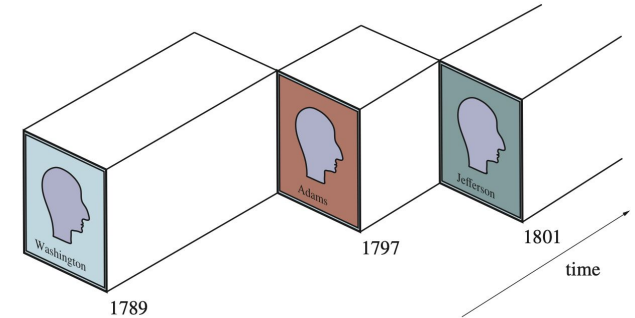
Reign of Elizabeth II immediately followed that of George VI, and the Reign of Elvis overlapped with the 1950s

Meets(ReignOf (GeorgeVI), ReignOf (ElizabethII))
 Overlap(Fifties, ReignOf (Elvis))
 Begin(Fifties) = Begin(AD1950)
 End(Fifties) = End(AD1959)

Fluents and Objects

- Objects can be viewed as generalized events

- A chunk of space-time
- USA is an event that started in 1776 with 13 states
- Population(USA) will be

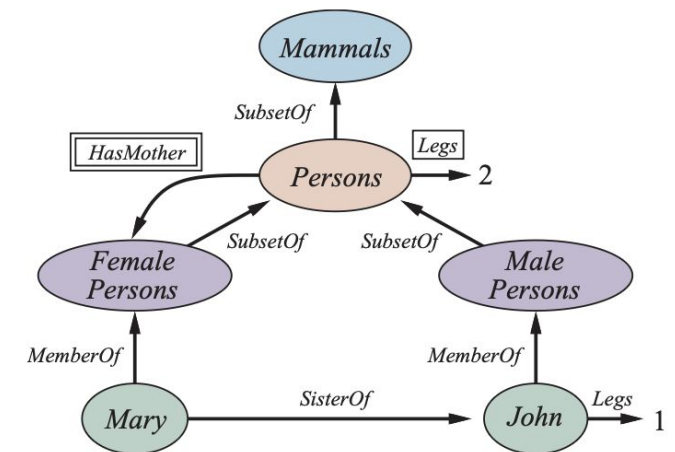


- Can President(USA) be a term if time is involved?

- President(USA) can denote a single object that consists of different people along with time intervals.
- T (Equals(President(USA), GeorgeWashington), Begin(AD1790), End(AD1790))
 - Here we use Equals instead of = because = cannot have predicate as an argument
 - Also, GeorgeWashington and President(USA) are not interpreted as same objects

Reasoning Systems with Category

- Two commonly used systems of organizing and reasoning about categories:
 - Semantic Networks
 - Description Logics
- **Semantic Networks:** Graph representation of categories, objects and relationships among them
 - **HasMother** is a property
 - How many legs does Mary have?
 - How many legs does John have?
- Reasoning is through graph traversal
 - Simple and efficient



Semantic Networks

- Multiple inheritance can lead to contradicting inferences

- We will discuss this later

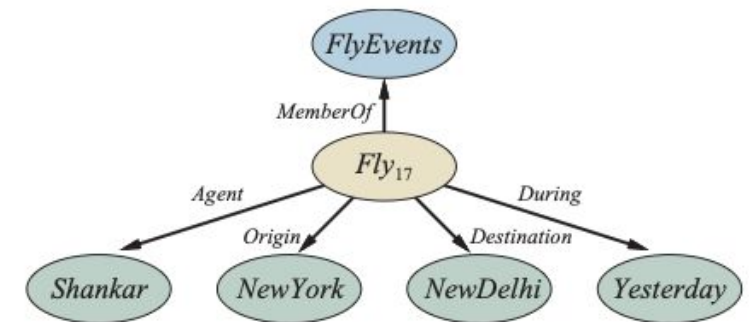
- How to assert Shankar Flew from NewYork to New Delhi Yesterday ?

- Strictly less expressive than first order logic

- In particular no existential quantifiers

- There are extensions

- At the cost of simplicity



Description Logics

- Evolved from semantic networks

- Mainly used to do the following inferences

- **Subumption** : A category is a subclass of another category
- **Classification** : An object belongs to a category
- **Consistency** : An object satisfies the property of the category to which it belongs

Concept → **Thing** | *ConceptName*
 | **And**(*Concept*,...) |
 | **All**(*RoleName*, *Concept*) |
 | **AtLeast**(*Integer*, *RoleName*) |
 | **AtMost**(*Integer*, *RoleName*) |
 | **Fills**(*RoleName*, *IndividualName*,...) |
 | **SameAs**(*Path*, *Path*) |
 | **OneOf**(*IndividualName*,...)

Path → [*RoleName*,...]

ConceptName → *Adult* | *Female* | *Male* | ...

RoleName → *Spouse* | *Daughter* | *Son* | ...

- Bachelor = And(Unmarried, Adult, Male)

- FemaleCSEProfs = AND(Faculty,

- All Men with :

- at least 3 sons, who are all unemployed, all of whom are married to doctors
- At most 2 daughters all of whom are professors in Math of Physics
- AND(Male, Atleast(3,Son), Atmost(2, Daughter),
 All(Son, AND(Unemployed, Married, All(Spouse,Doctor)),
 All(Daughter, AND(Professor, Fills(Department, Physics, Math))))

Description Logics

- Inference can be done in Polynomial time
- Description Logic does not have negation and only has limited disjunctions (in Fills and OneOf)

Concept → **Thing** | *ConceptName*
| **And**(*Concept*,...) | **All**(*RoleName*,*Concept*)
| **AtLeast**(*Integer*,*RoleName*) | **AtMost**(*Integer*,*RoleName*)
| **Fills**(*RoleName*,*IndividualName*,...) | **SameAs**(*Path*,*Path*)
| **OneOf**(*IndividualName*,...)

Path → [*RoleName*,...]

ConceptName → *Adult* | *Female* | *Male* | ...

RoleName → *Spouse* | *Daughter* | *Son* | ...

Mental models

- Till now agents perceive things from which they form their belief and derive new beliefs
 - They do not have beliefs about their own beliefs (Introspection)
- Scenario 1:
 - Alice : What is the square root of 400?
 - Bob : I do not know
 - Alice :
Bob realizes that if he puts more more thought, he should be able to figure out the answer. Think harder.
- Scenario 2:
 - Alice : Is Anantha's desktop in his office switched on or switched off ?
 - Bob : I do not know
 - Alice :
Bob should be able to say, there is no way he can get the answer by thinking harder. Think harder.
- Scenario 2:
 - Alice : Is Anantha's desktop in his office switched on or switched off ?
 - Bob : I do not know
 - Alice :
Bob : Does Anantha know?
Bob : Yes, Anantha obviously! know?
Bob should be able to reason about Anantha's knowledge.
- For such reasoning, the agents should be able to model the knowledge of other agents

Knowledge

- **Knows** : can be a predicate
 - $\text{Knows}(\text{Alice}, \text{Fly}(\text{Superman}))$ here $\text{Fly}(\text{Superman})$ is a predicate and not a term
 - Since $\text{Clark} = \text{Superman}$ we have :
 $\{ \text{Clark} = \text{Superman}, \text{Knows}(\text{Alice}, \text{Fly}(\text{Superman})) \} \models \text{Knows}(\text{Alice}, \text{Fly}(\text{Clark}))$
- **Referential Transparency** : Property depends only on the object that the term is referring to, not the syntax of the term itself
 - If $2+2 = 4$ and $4 < 5$ then $2+2 < 5$
- When modelling Knowledge of agents, we want Referential Opacity
 - Modal Logic can do this

Modal Logic

- Modal operators take sentences as inputs
 - A knows P is represented as $K_A P$
- Syntax is same as First Order Logic with additional operator $K_A P$
- Example (with propositions) : Alice is in Chennai and Bob is in Bangalore.
 - Alice knows that it is raining in Chennai
 $\blacksquare K_{\text{Alice}} (\text{RainingInChennai})$
 - Bob does not know whether it is raining in Chennai
 $\blacksquare \neg K_{\text{Bob}} (\text{RainingInChennai}) \wedge \neg K_{\text{Bob}} (\neg \text{RainingInChennai})$
 - Bob knows that it is either raining in Chennai or not
 $\blacksquare K_{\text{Bob}} (\text{RainingInChennai} \vee \neg \text{RainingInChennai})$
 - Bob knows that Alice knows whether it is raining in Chennai
 $\blacksquare K_{\text{Bob}} (K_{\text{Alice}} (\text{RainingInChennai}) \vee K_{\text{Alice}} (\neg \text{RainingInChennai}))$

Modal Logic

- Models in Modal Logic consists of Possible worlds
 - There is an indistinguishability relationship between the worlds for a given agent based on whether the agent can distinguish between them (also called accessibility relation).
 - $K_A P$ is true at a world w if P is true in all worlds that are accessible by A from w .
- Properties of Knowledge:
 - $K_A (\alpha \Rightarrow \beta) \Rightarrow (K_A \alpha \Rightarrow K_A \beta)$ (Agent knows Modus Ponens)
 - $K_A \alpha \Rightarrow \alpha$ (Agent's knows cannot be false)
 - $K_A \alpha \Rightarrow K_A K_A \alpha$ (Agent can positively introspect)
 - $\neg K_A \alpha \Rightarrow K_A \neg K_A \alpha$ (Agent can negatively introspect)
- Above properties + Axioms of Propositional Logic is a sound and complete axiom system for Propositional Epistemic Logic
- Logical Omniscience : Agent knows all consequences of the axioms $\alpha \Rightarrow K_A \alpha$
 - True for ideal agents
 - Real agents are assumed to have resource bounded inference power

Modal Logic

- Knowledge over predicates.

- Can describe subtle properties

- Example:

- Alice knows that someone killed Mary

- $K_{\text{Alice}} (\exists x \text{ Killed}(x, \text{Mary}))$

- Alice knows who killed Mary

- $\exists x K_{\text{Alice}} (\text{Killed}(x, \text{Mary}))$

- There are other Modal Logics variants:

- **Temporal Logic** : To reason about Time Always / Until / Sometime in Future / ...
- **Doxastic Logic** : Belief (need not be true)
- **Deontic Logic** : It is obligatory for adults to pay taxes / Request / Permission
- **Term Modal Logic** : There exists an Agent who knows alpha

Reasoning about exceptions

- In a semantic network, a subcategory can override a property
- Humans believe in default reasoning unless we are presented with an opposing evidence.
 - If my friend says I have parked my car outside, I will always assume that it has 4 wheels.
 - If I see someone limping, I will always assume that his leg is hurt.
- These are examples of nonmonotonic situations
 - Set of inferences does not monotonically grow as new evidence is presented
- There are two logics to deal with such scenarios:
 - Circumscription
 - Default logic

Circumscription

- Circumscribed Predicates are false for every object except for those specified otherwise
 - $\text{Bird}(x) \wedge \neg \text{Abnormal}_1(x) \Rightarrow \text{Flies}(x)$
- Here Abnormal_1 is a circumscribed predicate
 - It is assumed to be false for every object unless specified otherwise
 - We can add $\text{Penguin}(x) \Rightarrow \text{Abnormal}_1(x)$
- Handling multiple inheritance:
 - Let Alice is a scientist (and hence believes Bigbang created the universe)
She is also a believer of god (and hence believes God created the universe and not bigbang).
 - $\text{Scientist}(\text{Alice}) \wedge \text{Theist}(\text{Alice})$
 - $\text{Scientist}(x) \wedge \neg \text{Abnormal}_2(x) \Rightarrow \text{BelievesBigBang}(x)$
 - $\text{Theist}(x) \wedge \neg \text{Abnormal}_3(x) \Rightarrow \neg \text{BelievesBigBang}(x)$
 -
 - If $\text{Abnormal}_3(\text{Alice})$ is in the knowledgebase, then we can conclude that she believes in god but also in Bigbang
 - But if we have both $\text{Abnormal}_2(\text{Alice})$ and $\text{Abnormal}_3(\text{Alice})$?
 - Inferencing algorithm (rightly) concludes that it cannot say either.

Default Logic

- Example 1:
 - $\text{Bird}(x) : \text{Flies}(x) / \text{Flies}(x)$
 - If $\text{Bird}(x)$ is true and if $\text{Flies}(x)$ is consistent with the KB then $\text{Flies}(x)$ is concluded by default
- General Rule looks like : $P : J_1 \dots J_n / C$
 - If P holds and $J_1 \dots J_n$ is consistent with the KB then conclude C
- Example 2:
 - $\text{Scientist}(x) : \text{BelievesBigBang}(x) / \text{BelievesBigBang}(x)$
 - $\text{Theist}(x) : \neg \text{BelievesBigBang}(x) / \neg \text{BelievesBigBang}(x)$
 - $\text{Scientist}(\text{Alice}) \wedge \text{Theist}(\text{Alice})$
- In default logic, we conclude a property only if it holds in every maximal consistent extension of the initial set S

Nonmonotonic Logic

- Nonmonotonic logics are non-modular
 - Every exception needs to be handled in its own way
 - There is no common way to handle all rules at a time
- Should we have default for every rule?
 - Which rules should have default and which should not?
- My car brakes are always OK
 - Given no other information, the probability that my car's brakes are OK is high enough for me to drive without checking them
 - We can use probability / threshold also to deal with defaults

Belief Revision

- **Default** is the **default status**.
 - But **when new evidence is presented**, some birds do not fly anymore
 - So we need to remove agent's beliefs and all its consequences
- This process is called **Belief Revision**
 - **RETRACT(KB, P)**
 - Should remove P and also all inferences that were added to KB that were inferred using P
- **Example :**
 - $KB = \{ P \Rightarrow Q, Q, P \}$ (where Q was an inference) then $RETRACT(KB, P) = \{ P \Rightarrow Q \}$
 - $KB = \{ P \Rightarrow Q, Q, P, R \Rightarrow Q, R \}$ (where Q was an inference) then
 $RETRACT(KB, P) = \{ P \Rightarrow Q, R, Q \}$
- **Truth Maintenance Systems** are designed to do such revisions efficiently.

Truth Maintenance Systems

- Keep track of the order in which the sentences are added to $KB : P_1 \dots P_n$
 - To remove P_i first revert to the state just before P_i was added. Then add $P_{i+1} \dots P_n$
 - Costly if revision happens frequently
- Justification based Truth Maintenance : Keep track of justification for the inferences.
 - $KB = \{ P \Rightarrow Q, P \}$ then $Justification(Q) = \{ \{P \Rightarrow Q, P\} \}$
 - $KB = \{ P \Rightarrow Q, P, R \Rightarrow Q, R \}$ then $ustification(Q) = \{ \{P \Rightarrow Q, P\}, \{R \Rightarrow Q, R\} \}$
 - RETRACT(P) will remove every Justification set that contains P.
 - Time for RETRACT(P) is proportional to the number of sentences that were inferred using P.
 - When an inferred sentence is removed from the KB it is marked as out (but its justification is remembered)
 - If later one of the justification set is added back then we do not have to recompute the entire proof
- Assumption based Truth Maintenance : For each sentence keeps track of which assumptions would make the sentence true.

Inference in the age of LLMs

- LLMs are great but
 - They are wrong in many cases
 - We do not know when it is right and when it is wrong
- Can we make LLMs use some Knowledge base?
 - Inferencing takes a lot of time
 - What should be stored in the knowledge base?
- Current research
 - Inference guided LLMs : Get an output from LLM, ask how it arrived at the conclusion. Then use a logical system to verify that the out is correct
 - It is easy to check a proof than to come up with a proof.
 - Guide it by saying at this step the reasoning is incorrect, if there is some flaw.
 - When an LLM says “This is how I arrived at a conclusion” did it really do it or just saying the most likely thing that you will accept?