







Classification of states " Accessi bility" A state j is said to be accessible from a state i if Juso S.t. P. 20 If j is accessible from i, we write i-j i→j =)] a directed path from i to j in the transition diagram Communication: States i and j communicate if inj ad jni Denote communication by "=> Clain: (i) i c> i (reflexive) (Ti) if i co i then i co i (symmetric) (iii) if icos, jcot, then icok (transitive) PF: (i) & (ii) hold by definition For (iii) i - j, j -> k)]n,m st. P(n) >0 & P(m) >0



















Claim: 1 1633, i su positive recurrent

$$2334$$
 positive recurrent
 2334 positive recurrent
 2334

Period of state 1=2 Visit instants of state 2= {1,3, --- } Periol of state 2= 2 1 DE DE 2 - - NF N (\mathfrak{d}) Period of 0 = 1 = Period of N absorbing states Pyro Recall T: = nin { n70 (Xn=i3, iES Let "i be a recurrent state. Let "d" be the langest positive integer such that $\mathcal{Z} P(\mathcal{T}_{i} = kd) = 1.$ If d=1, then state i is aperiodic. d>1, then state i is periodic with period d. J. Equivalently if''' is a recurrent state with period d, $then <math>p_{i,i}^{(n)} = 0$ is in that are not positive integer mattiples of d. Periodicity is a class property! Claim! It ic) 3, then i and 3 have the same period.

If: Since
$$i \in 3$$
, $\exists n,m$ st. $p_{i,j}^{(n)} > 0$, $p_{i,j}^{(m)} > 0$
Hiso, $p_{i,j}^{(r+n+m)} \ge p_{i,j}^{(n)} p_{i,j}^{(m)} - (\omega)$
Suppose "i" has a period "d".
With $r=0$ in (ω), we obtain
 $p_{i,j}^{(n+m)} \ge p_{i,j}^{(n)} p_{i,j}^{(m)} > 0$
 $p_{i,j}^{(n+m)} \ge p_{i,j}^{(n)} p_{i,j}^{(m)} > 0$
 $p_{i,j}^{(r+m)} \ge 0 \Longrightarrow p_{i,j}^{(n)} p_{i,j}^{(n)} = 0$
 $p_{i,j}^{(r+n+m)} = 0 \Longrightarrow p_{i,j}^{(n)} p_{i,j}^{(n)} = 0$
 $p_{i,j}^{(r)} = 0$ for τ that isn't divisible by d.
 $\sum p_{i,j}^{(r)} = 0$ for τ that isn't divisible by d.
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Petermining recurrence in a finite state DIMC We will establish that (I) all states in a finite closed communicating class are positive recurrent (II) all states in a non-closed communicating class are transient From I 4 II : AU communications clones of a finite DIMC are finite, which implies that they are either possitive recurrent or transient. €03, EN3, E1, ___, N-1] clore d positive recurrent transient There are no "null-recurrent" communicating closes in a finite DTMC. Proof of (I): Let C be a finite closed communicating class. Then, See next page

Proof of claim II: Let C be a non-closed communicating class. FIEC jEC, such that i is not accessible from j, and Pi, 20 -> If the DTMC visita ; starting from i, it never returns to i $P(\widetilde{T}_{i} = \infty | \chi_{o} = i) \geq P_{i,i} \geq 0$ fecal W:= P(F; c ~ (x=i) So, $I - G_i = P(\widetilde{T}_i = \infty (\pi_0 = i) > 0$ \Rightarrow $\mathcal{L}_{i} < 1$ =) i is transient =) C is transfert

Suppose i \$ j. Irore duciblity =) j = i =) j-) î =) p(m) >0 for some $\begin{array}{c|c} (n) & & (m) & (n-m) \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & i \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & i \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & i \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & i \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta}_i i & \dot{\delta} \\ \hline \dot{\delta}_i \dot{\delta} & & \dot{\delta} \\ \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} \\ \dot{\delta} & \dot{\delta} & \dot{\delta} \\ \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} \\ \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} \\ \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} \\ \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} \\ \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} \\ \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} \\ \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} & \dot{\delta} \\ \dot{\delta} & \dot{$ Anzm Since $p_{j,j}^{(n)} \longrightarrow D$ & $p_{j,j}^{(m)} = 70$, we obtain p(n-m) -> 0 as (n-m) -> 00 Coming next: (I) $\lim_{n \to \infty} P(X_n = j) = 0$ $\forall j$ (II) Suppose the state space is finite. Then, I at least one recurrent state. (II) Suppose the state space is firste and the DTMC is irreducible. Then, all states are recurrent.

Lecture - 20

Pf of (I): Recall we have an irreducible transient DTMC. Let a= (a;); es be the instal distribution $P(X_n=j) = \sum_{i \in i} \alpha_i P_{i,j}^{(n)}$ $\lim_{n \to \infty} P(\chi_n = i) = \lim_{n \to \infty} \sum_{i \in I} \alpha_i P_{i,i}^{(n)}$ Suppose ISI < 00. Then, we can interchage the limit f sum $\lim_{n \to \infty} P(\chi_{n=j}) = \sum_{i \in S} a_i \left(\lim_{n \to \infty} P_{i,j}^{(n)} \right)$ $\lim P(x_1=j) = 0$ <u>v ~) ∞</u> Suppose Ist=00. Then, ve can use dominated convergence theorem (OCT) to justify the interchange of lim & sum. A precise argument is given below. Let F be a finite subset of S. $\sum_{i \in S} a_i P_{i,i}^{(n)} = \sum_{i \in F} a_i P_{i,i}^{(n)} + \sum_{i \notin F} a_i P_{i,i}^{(n)}$ $\leq \sum_{i \in C} \alpha_i \rho_{ij}^{(n)} + \sum_{i \notin F} \alpha_i$ 1.300 5 Q; FAS, D

So, $\lim P(x_n = \delta) = 0$. n→∞ (II) Suppore the state space is finite. Then, it at least one recurrent state. $\frac{Pf:}{\sum_{j \in S} P_{i,j}^{(n)} = 1}$ Assume every state is transient, i.e., This leads to a contradiction. (II) Suppose the state space & firste and the DIMC is irreducible. Then, all states are recurrent. Pf. From part (I), we have at least one recurrent state. (7) DTMC is irreducible =) All states are recurrent. Coming later: Ina faite îrreducible DTMC, all states are positive recurrent.

A bunch of claims: For an irreducible, recurrent Dim, with e as constructed above, we have (i) $e_{k} = 1$, (ii) $\sum_{i \neq s} e_{i} = m_{k}$, (iii) $e_{i} = e_{i}$ $(iv) \quad 0 < e_i < \infty \quad \forall i \in S.$ Proof - (i) Recall $W_i = \sum_{m=1}^{T_E} I(x_m=i) \otimes P_i = E(W_i | x_s=k)$ So, WE=1& PE=1. (ii) $T_k = \sum_{k \in S} W_i$ $m_{k} = E(T_{k}|X_{0}=k)$ $= E\left(\sum_{i \in V} W_i \mid X_{o} \geq k\right)$ monotone convergence $y = \sum E(W, | X_0 = E)$ the original of (See Sec S. 6 in sbook, (See Sec Stirzater there) Grinnett- ex (133) Grinnett-1 = $\leq e_i$ (Tii) To show P=PP $e_{j} = E\left(\sum_{m=1}^{j_{k}} I(x_{m}=j) | X_{o}=k\right)$ $Let E_{k}(\cdot) = E(\cdot (\chi_{s-k}) + P_{k-1} + P(\cdot | \chi_{s-k}))$

So,
$$e_{3} = E_{k} \left(\sum_{m=1}^{k} I\left(X_{m}=\hat{\delta}\right) \right)$$

$$= E_{k} \left(\sum_{m=1}^{k} I\left(X_{m}=\hat{\delta}, T_{k}, \mathbb{Z}^{m}\right) \right)$$

$$= \sum_{m=1}^{\infty} P_{k} \left(X_{m}=\hat{\delta}, T_{k}, \mathbb{Z}^{m}\right)$$
Notice that
 $P_{k} \left(X_{m}=\hat{\delta}, T_{k}, \mathbb{Z}^{m}\right)$

$$= \sum_{n=1}^{\infty} P_{k} \left(X_{m-1}=i, X_{m}=\hat{\delta}, T_{k}, \mathbb{Z}^{m}\right)$$

$$= \sum_{n=1}^{\infty} P_{k} \left(X_{m-1}=i, X_{m}=\hat{\delta}, T_{k}, \mathbb{Z}^{m}\right) P_{k} \left(X_{m-1}=i, T_{k}, \mathbb{Z}^{m}\right)$$
we show mater graphs
 $\overline{P} = \sum_{n=1}^{\infty} P_{i} \left(X_{m}=\hat{\delta}, X_{m-1}=i, T_{k}, \mathbb{Z}^{m}\right) P_{k} \left(X_{m-1}=i, T_{k}, \mathbb{Z}^{m}\right)$
Notice that
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Notice the store graphs
 $\overline{P} = \sum_{n=1}^{\infty} P_{i} \left(X_{m}=\hat{\delta}, X_{m-1}=i, T_{k}, \mathbb{Z}^{m}\right) P_{k} \left(X_{m-1}=i, T_{k}, \mathbb{Z}^{m}\right)$
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 $\overline{P} = \sum_{n=1}^{\infty} P_{i} \left(X_{m-1}=i, T_{m-1}=i, T_{m}, T_{m}, T_{m}=i, T_{m}, T_{m}=i, T$

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Two (random) Stopping time.

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 of detour)
Back to proof of $f = \ell P$
We have,
 $P_{\ell}(X_{v} = \delta), T_{k} \ge m) \ge \sum_{i \in S} P_{i,i} P_{k}(X_{m-1} = i, T_{k} \ge m)$
 $Vurly hill in the expression for P_{δ} , we obtain
 $P_{\delta} = \sum_{m=1}^{\infty} \sum_{i \in S} P_{i,i} P_{k}(X_{m-1} = i, T_{k} \ge m)$
 $= \sum_{i \in S} \sum_{m=1}^{\infty} P_{i,i} P_{k}(X_{m-1} = i, T_{k} \ge m)$
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 $= \sum_{i \in S} P_{i,i} \sum_{m=1}^{\infty} P_{k}(X_{m-1} = i, T_{k} \ge m)$
 $= \sum_{i \in S} P_{i,i} \sum_{m=1}^{\infty} P_{k}(X_{k} \ge i, T_{k} \ge T(1))$
 $= \sum_{i \in S} P_{i,i} \sum_{m=1}^{\infty} P_{k}(\sum_{m=1}^{\infty} I \le X_{w} = i, T_{k} \ge T(1))$
 $= \sum_{i \in S} P_{i,i} \sum_{i \in S} P_{i,i} E_{k}(\sum_{m=0}^{\infty} I \le X_{w} = i)$
 $P_{i} = \sum_{i \in S} P_{i,i} E_{i,i} E_{k}(\sum_{m=0}^{\infty} I \le X_{w} = i)$$

Big facts about stationary distributions Irreducible DTMC. (not nocenanily) (i) if some state is positive recurrent, then Ja stationary distribution. PE: Suppose k is positive recurrent $_{2}) m_{E} < \infty$ From earlier thm, we have a stationcry measure lie, l= lP Sconstructed uping state K Set $T_r = \frac{P_i}{m_k}$ then TT = TT P & also TT is a distribution Since, $\Xi T := \frac{\Xi e_i}{m_E} = 1$ (ii) If I a stationary distribution TT, then (a) each state is positive recurrent

$$T_{ij} = \sum_{i=1}^{n} T_{ij} \frac{p(n)}{n - 900} = 0 \quad (for preficient loss is the set in the set in$$

Pf: Uses "Coupling". Let Stag be an independent copy of the given DTMC, i.e., {Y,3 how the same state, say S, Space & t.p.m.Pas &x.3. $Z_n = (X_n, Y_n)$ Is Zn a DTMC? Yes. State space SXS $f.p.m. for Z_n:$ $P = P(Z_n = (k, l) | Z_n = (i, j))$ $P_{ij,kl}$ $\sum_{i=1}^{n} \frac{2}{i} P(X_{n+i} = k | X_n = i) P(Y_{n+i} = L|Y_n = i)$ $= P_{i,k} P_{i,k}$ Is 2 irreducible? Yes. ZX_3, 54_3 to îrreducible, operiodic (Given) JN s.t. <u>YnZN</u> p(n) p(n) >0 <u>YnZN</u> h.w. p(n) p(n) >0 <u>YnZN</u> aperiodicity. Does S2n3 have a stationary distribution? Yes. Let v= (V;; i,jESxS) denote the Stationary Listribution of §2n3.

Conditional on STEn3, both Xn & Yn have the same distributions. E Strong Markov property If two indep. DIMCA with identical t. p. m. start in the same state, say Xo=5, then The Istribution of X. 4 Th are going to be strong Marten the Same A rimiter observation holds for a strong Marten Storping time T, owhen one looks at the event (TEng) $P_{i,k}^{(n)} = P_i(X_n = k)$ = $P_{i}(X_{n}=k, T\leq n) + P_{i}(X_{n}=k, T>n)$ with the hore $D = P.(Y_n = k, T \leq n) + P;(X_n = k, T \geq n)$ the same same sign bulkon $Y_n = k, T \leq n + P;(X_n = k, T \geq n)$ with 875n3, $\leq P_{1,2}(Y_n=k) + P_{1,2}(T_n)$ $= P_{i,k}^{(n)} + P_{i,k}^{(T7n)}$ So, $p_{i,k}^{(n)} - p_{i,k}^{(n)} \leq P_{i\delta}$ (T>n) Swap it of f orpeat the argument above to get $P_{i,k}^{(n)} - P_{i,k}^{(n)} \leq P_{i,j}(T > n)$ $|P_{i,k}^{(n)} - P_{\delta,k}^{(n)}| \leq P_{i\delta}(T > n)$ 0۴,

As n-2 00, P; (T>n) -> O because of (70) So, $p(n) - p(n) \rightarrow 0$ as $n \rightarrow \infty$. So, it lim Pik exists, then it does not noo it depend on starting state i lim Pite exists because TT_k - p⁽ⁿ⁾ k k $= \sum_{i=1}^{n} \prod_{i=1}^{n} \left(P_{i,k}^{(n)} - P_{i,k}^{(n)} \right) \qquad \text{since } \Pi = \Pi P P S \Pi_{i=1}^{n}$ n-200) O (by a version of dominated convergence theorem. see p. 30 for interne). $p \xrightarrow{(n)} \longrightarrow TT \xrightarrow{(n)} \infty$ So,

n mo mo a.s. under P(. /X2i) The claim follows. Some examples: $a + \frac{1}{2} + \frac{1}{2}$ 2+62 (1) Two state DTMC: $\pi = \pi P$ (=) $\pi = 2\pi + (1-\beta)\pi_2$ $\Pi_{1} \simeq (1 \sim 1) \Pi_{1} + \beta \Pi_{2}$ $\Pi_1 + \Pi_2 = 1$ If you solve, then $T_{1} = \frac{1-\beta}{2-2-\beta}$, $T_{2} = \frac{1-2}{2-2-\beta}$ Check the finding in Example 1 on p.2 of this nate. H.W. Find stationary TT for the periodic OTMC in Example 2 on p. 3. of this note.

In this case, the DTMC is positive recurrent.
H.W.: If
$$S(t;=\infty)$$
, determine null recurrent.
transience of "success runs" PTMC.
Special case: $p_{1}=p$ $\forall i$
 $f_{1}=p^{2}=p$ $\forall i$
 $T_{1}=p^{2}(1-p) \leftarrow Treemetric detections
Simplified version of hoogle lagerant
Suppose there are 4 webpages
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Question! Now to compute long run avg. cost objective? long. run angunt $\mathcal{I}(i) = \lim_{n \to \infty} \frac{1}{n} \sum_{m=1}^{n} E(c(X_m) | X_n = i)$ (*) Suppose DTMC is irreducible, positive recurrent (2) 7 a stationary distribution Ti) What we will show is $T(i) = J \forall i$. Claim: For a DTMC satisfing (7), with $\sum_{j=1}^{n} T_{j} [c(j)] < \infty$, we have J(i) = $\sum_{j \in S} T_j C(j)$ avg. cost Same for all that states. $5(i) = \lim_{n \to \infty} \frac{1}{n} \sum_{m=1}^{n} E(c(X_m) | X_n = i)$ Rf : $= \lim_{n \to \infty} \frac{1}{n} \sum_{j \in S} \frac{V_{i,j}^{(n)} c(j)}{j \in S},$ where $V_{i,j}^{(n)} = expected # visits to j in 21 --- n3$ Starting in state i $\mathcal{J}(i) = \lim_{n \to \infty} \frac{\sum V_{i,j}}{\sum S \in S} \frac{\nabla_{i,j}}{n} c(i)$ So, = $\sum_{i \in S} \left(\underset{n > \infty}{\underset{n > \ldots}{\underset{n > \ldots}{n > \ldots}{n > \ldots}{\underset{n > \ldots}{n > \ldots}{n > \ldots}{n > \ldots}{n > \ldots}$

J(i) = <u>S</u> TT, C(j). Le we use de possitive recurrence + JES TT, C(j). Le we use de possitive recurrence + GES Gesoro convergence (see big fact I) For allerage cost objective, a good reforence les "Dynamic programming 4 optimal control, Vol. I" by D. P. Bertsetas, Athena Scientific. Example: Brand Switching (-3. $P = \begin{array}{c} A & B \\ A & 0.1 \\ 0.2 & 0.7 \\ 0.2 & 0.4 \\ 0.4 & 0.4 \\ 0.1 & 0.3 & 0.6 \end{array}$ P(buying B | given bought A Lost time) - 0.2 Costr: c(A) = 6, c(B) = 5, c(C) = 4. Finite irreducible DIMC =) $\exists T S.f. TT = TTP$, $\Xi TT = 1$ $\Pi_{\mu} = 0.13$, $\Pi_{\beta} = 0.32$, $\Pi_{c} = 0.55$ Long vun average cost J = TTA ((A) + TTB c (B) + TTC c (C) = 4.58 So, the Customer spinds 4.58 Custatever currency on each vicit in the long ran. (End of DTMCs>