

Shape Analysis

Rupesh Nasre.

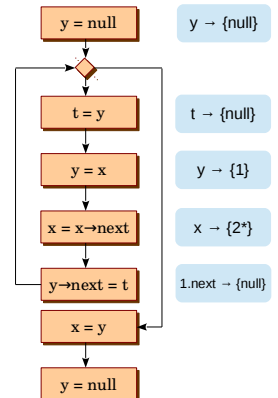
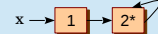
CS6843 Program Analysis
IIT Madras
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Limitations of Pointer Analysis

```
listReverse(List x) {
  assert("x is an acyclic singly linked list");

  for (y = null; x;) {
    t = y;
    y = x;
    x = x->next;
    y->next = t;
  }
  x = y;
  t = null;
  y = null;
}
```

We model the list as a two node structure.



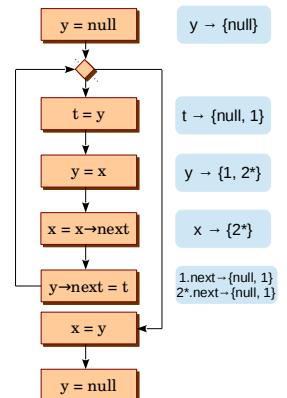
Outline

- Limitations of pointer analysis
- Identify lists
- Identify trees, DAGs, cyclic graphs
- Identifying rotations
- List reversal and other transformations

Limitations of Pointer Analysis

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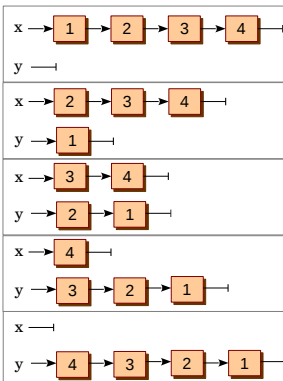


Limitations of Pointer Analysis

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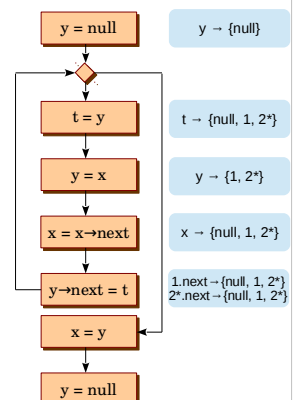
We want to check if x points to a singly linked list at the end of listReverse. That is, $x \rightarrow \{4\}$, $4.\text{next} \rightarrow \{3\}$, ..., $1.\text{next} \rightarrow \{\text{null}\}$



Limitations of Pointer Analysis

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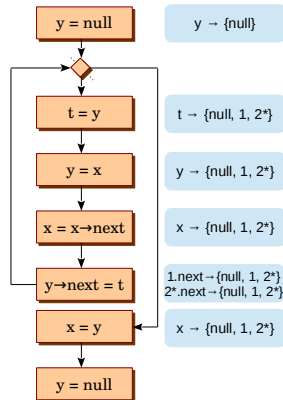
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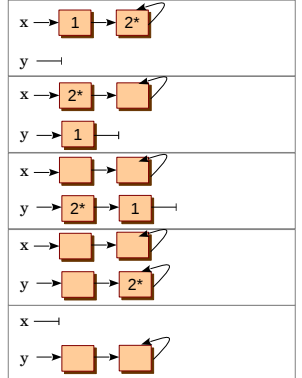
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Shape Analysis

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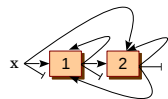


Maintain additional information with 2^* that it is acyclic. Use the fact that node removal maintains acyclicity.

Limitations of Pointer Analysis

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    y->next = t;
  }
  x = y;
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}
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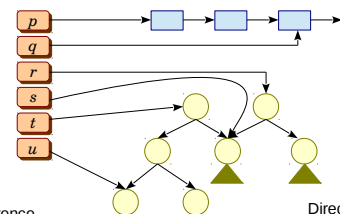
Tree, DAG, Cycle?

- Maintains three data structures:
 - Interference matrix: encodes common reachability
 - Direction matrix: encodes direct reachability
 - Shape
- Performs iterative data-flow analysis to update D, I and shape information

Shape Analysis

- Identify structural / topological properties of a data structure under manipulation.
- Usually categorized as slist, tree, DAG or cycle.
- Precision reduces along slist $\rightarrow$ tree $\rightarrow$ DAG $\rightarrow$ cycle.

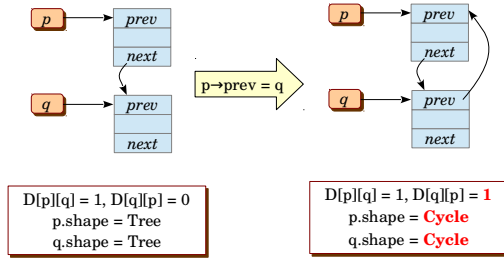
Tree, DAG, Cycle?



	p	q	r	s	t	u
p	1	1	0	0	0	0
q		1	0	0	0	0
r			1	1	1	0
s				1	1	0
t					1	1
u						1

	p	q	r	s	t	u
p	1	1	0	0	0	0
q	0	1	0	0	0	0
r	0	0	1	1	1	0
s	0	0	0	1	1	0
t	0	0	0	0	1	1
u	0	0	0	0	0	1

Shape Estimation



Inference Rules

$p = \text{malloc}(\dots)$	D_kill and I_kill sets same as for allocation statement.
$p = q$	$D_gen = \{D[s][p] \mid D[s][q] \text{ and } s \neq p\} \cup$ $\{D[p][s] \mid D[q][s] \text{ and } s \neq p\} \cup$ $\{D[p][p] \mid D[q][q]\}$
$p = q \rightarrow f$	$I_gen = \{I[p][s] \mid I[q][s] \text{ and } s \neq p\} \cup$ $\{I[p][p] \mid I[q][q]\}$
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = \text{null}$	
$p \rightarrow f = q$	$p.shape = q.shape$
$p \rightarrow f = \text{null}$	

Inference Rules

$p = \text{malloc}(\dots)$	
$p = q$	
$p = q \rightarrow f$	
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = \text{null}$	
$p \rightarrow f = q$	
$p \rightarrow f = \text{null}$	

Inference Rules

$p = \text{malloc}(\dots)$	D_kill and I_kill sets same as for allocation statement.
$p = q$	$D_gen = \{D[s][p] \mid I[s][q] \text{ and } s \neq p\} \cup$ $\{D[p][s] \mid D[q][s] \text{ and } s \neq p \text{ and } s \neq q\} \cup$ $\{D[p][q] \mid q.shape == Cycle\} \cup$ $\{D[p][p] \mid D[q][q]\}$
$p = q \rightarrow f$	$I_gen = \{I[p][s] \mid I[q][s] \text{ and } s \neq p\} \cup$ $\{I[p][p] \mid I[q][q]\}$
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = \text{null}$	
$p \rightarrow f = q$	$p.shape = q.shape$
$p \rightarrow f = \text{null}$	

Inference Rules

$p = \text{malloc}(\dots)$	$D_kill = \{D[p][s] \mid D[p][s] == 1\} \cup \{D[s][p] \mid D[s][p] == 1\}$ $I_kill = \{I[p][s] \mid I[p][s] == 1\}$
$p = q$	$D_gen = \{D[p][p]\}$ $I_gen = \{I[p][p]\}$
$p = q \rightarrow f$	
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = \text{null}$	$p.shape = Tree$
$p \rightarrow f = q$	
$p \rightarrow f = \text{null}$	

Inference Rules

$p = \text{malloc}(\dots)$	Processing is the same as for $p = q$ statement. This means the analysis loses field-sensitivity.
$p = q$	
$p = q \rightarrow f$	
$p = \&(q \rightarrow f)$	A recent work from IITK (Dasgupta, Karkare, Reddy) addresses this issue.
$p = q \text{ op } k$	
$p = \text{null}$	
$p \rightarrow f = q$	
$p \rightarrow f = \text{null}$	

Inference Rules

<code>p = malloc(...)</code>	D_kill and I_kill sets same as for allocation statement.
<code>p = q</code> <code>p = q→f</code> <code>p = &(q→f)</code> <code>p = q op k</code> <code>p = null</code>	D_gen = { } I_gen = { } p.shape = Tree
<code>p->f = q</code> <code>p->f = null</code>	

Example

```
listReverse(List x) {
  assert("x is an acyclic singly linked list");

  for (y = null; x;) {
    t = y;
    y = x;
    x = x→next;
    y→next = t;
  }
  x = y;
  t = null;
  y = null;
}
```

Interference			
	x	y	t
x		0	0
y			0
t			

Direction			
	x	y	t
x	1	0	0
y	0	0	0
t	0	0	0

Shape	
x	tree
y	tree
t	tree

Inference Rules

<code>p = malloc(...)</code>	D_kill = { }, I_kill = { }
<code>p = q</code> <code>p = q→f</code> <code>p = &(q→f)</code> <code>p = q op k</code> <code>p = null</code>	D_gen = {D[r][s] D[r][p] and D[q][s]} I_gen = {I[r][s] D[r][p] and I[q][s]} D[q][p] and D[s][q] ⇒ s.shape = Cycle D[q][p] and D[s][p] ⇒ s.shape = Cycle !D[q][p] and D[s][p] and I[s][q] and q.shape == Tree ⇒ s.shape = max(s.shape, DAG) !D[q][p] and D[s][p] and q.shape != Tree ⇒ s.shape = max(s.shape, q.shape)
<code>p->f = q</code> <code>p->f = null</code>	

	Tree	DAG	Cycle
Tree	Tree	DAG	Cycle
DAG	DAG	DAG	Cycle
Cycle	Cycle	Cycle	Cycle

max(shape1, shape2)

Example

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listReverse(List x) {
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  for (y = null; x;) {
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    y = x;
    x = x→next;
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  }
  x = y;
  t = null;
  y = null;
}
```

Interference			
	x	y	t
x		0	0
y			0
t			

Direction			
	x	y	t
x	1	0	0
y	0	0	0
t	0	0	0

Shape	
x	tree
y	tree
t	tree

Inference Rules

<code>p = malloc(...)</code>	D_kill = { }, I_kill = { }
<code>p = q</code> <code>p = q→f</code> <code>p = &(q→f)</code> <code>p = q op k</code> <code>p = null</code>	D_gen = { } I_gen = { } No changes to the shape of p.
<code>p->f = q</code> <code>p->f = null</code>	

Example

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Interference			
	x	y	t
x		0	0
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Direction			
	x	y	t
x	1	0	0
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Shape	
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y	tree
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}
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Interference			
	x	y	t
x		1	0
y			0
t			

Direction			
	x	y	t
x	1	1	0
y	1	1	0
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Interference			
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y			1
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x	1	0	0
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```

We need to assume a finite representation for the data structure.

Interference			
	x	y	t
x		1	0
y			0
t			

Direction			
	x	y	t
x	1	0	0
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Interference			
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	x	y	t
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```

Since we do not model fields, we can't say that x is unreachable from y.

Interference			
	x	y	t
x		1	0
y			0
t			

Direction			
	x	y	t
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Shape	
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Shape

x	cycle
y	cycle
t	cycle

Improvements

- SSA form
- Field-sensitivity
- Heap modeling

Summary

- Shape analysis helps several transforms.
- Existing techniques often trade off precision for efficiency.
- We are still far away from a precise and scalable analysis.