

Shape Analysis

Rupesh Nasre.

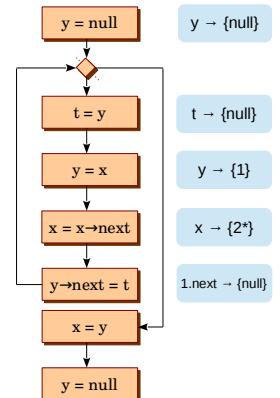
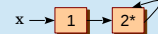
CS6843 Program Analysis
IIT Madras
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Limitations of Pointer Analysis

```
listReverse(List x) {
  assert("x is an acyclic singly linked list");

  for (y = null; x;) {
    t = y;
    y = x;
    x = x->next;
    y->next = t;
  }
  x = y;
  t = null;
  y = null;
}
```

We model the list as a two node structure.



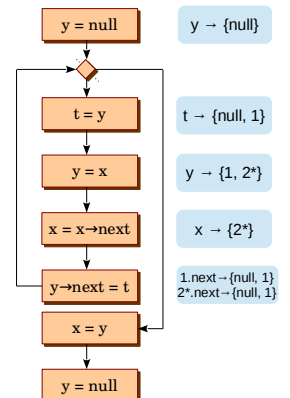
Outline

- Limitations of pointer analysis
- Identify lists
- Identify trees, DAGs, cyclic graphs
- Identifying rotations
- List reversal and other transformations

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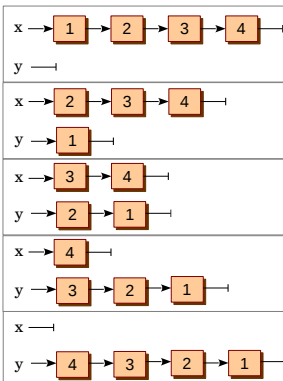


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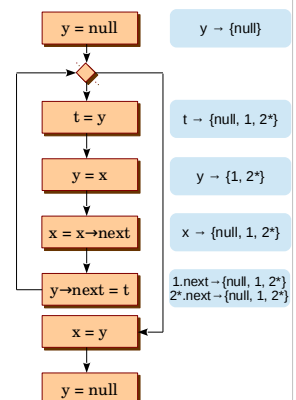
We want to check if x points to a singly linked list at the end of listReverse. That is, $x \rightarrow \{4\}$, $4.\text{next} \rightarrow \{3\}$, ..., $1.\text{next} \rightarrow \{\text{null}\}$



Limitations of Pointer Analysis

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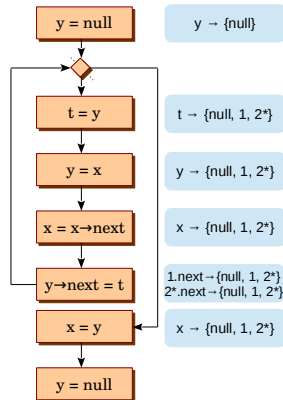
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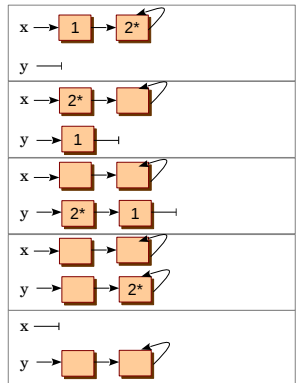
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Shape Analysis

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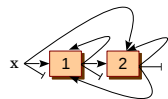


Maintain additional information with 2* that it is acyclic. Use the fact that node removal maintains acyclicity.

Limitations of Pointer Analysis

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```



Control flow graph for listReverse. The graph shows the state of variables at each point in the execution. The initial state is y = null. The loop body consists of t = y, y = x, x = x->next, and y->next = t. The final state is x = y and y = null.

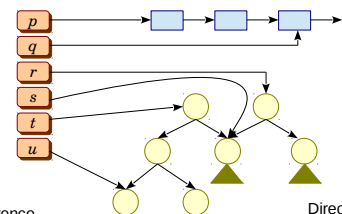
Tree, DAG, Cycle?

- Maintains three data structures:
 - Interference matrix: encodes common reachability
 - Direction matrix: encodes direct reachability
 - Shape
- Performs iterative data-flow analysis to update D, I and shape information

Shape Analysis

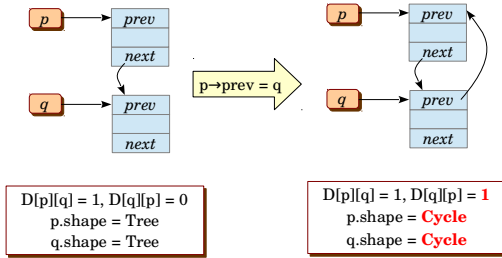
- Identify structural / topological properties of a data structure under manipulation.
- Usually categorized as slist, tree, DAG or cycle.
- Precision reduces along slist → tree → DAG → cycle.

Tree, DAG, Cycle?



Interference							Direction						
	p	q	r	s	t	u		p	q	r	s	t	u
p	1	1	0	0	0	0		1	1	0	0	0	0
q		1	0	0	0	0		0	1	0	0	0	0
r			1	1	1	0		0	0	1	1	0	0
s				1	1	0		0	0	0	1	0	0
t					1	1		0	0	0	1	1	1
u						1		0	0	0	0	0	1

Shape Estimation



Inference Rules

$p = \text{malloc}(\dots)$	D_kill and I_kill sets same as for allocation statement.
$p = q$	$D_gen = \{D[s][p] \mid D[s][q] \text{ and } s \neq p\} \cup \{D[p][s] \mid D[q][s] \text{ and } s \neq p\} \cup \{D[p][p] \mid D[q][q]\}$
$p = q \rightarrow f$	$I_gen = \{I[p][s] \mid I[q][s] \text{ and } s \neq p\} \cup \{I[p][p] \mid I[q][q]\}$
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = \text{null}$	
$p \rightarrow f = q$	$p.shape = q.shape$
$p \rightarrow f = \text{null}$	

Inference Rules

$p = \text{malloc}(\dots)$
 $p = q$
 $p = q \rightarrow f$
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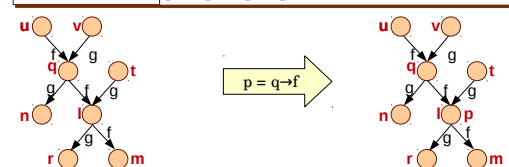
$p = \text{malloc}(\dots)$	D_kill and I_kill sets same as for allocation statement.
$p = q$	$D_gen = \{D[s][p] \mid I[s][q] \text{ and } s \neq p\} \cup \{D[p][s] \mid D[q][s] \text{ and } s \neq p \text{ and } s \neq q\} \cup \{D[p][q] \mid q.shape == Cycle\} \cup \{D[p][p] \mid D[q][q]\}$
$p = q \rightarrow f$	$I_gen = \{I[p][s] \mid I[q][s] \text{ and } s \neq p\} \cup \{I[p][p] \mid I[q][q]\}$
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = \text{null}$	
$p \rightarrow f = q$	$p.shape = q.shape$
$p \rightarrow f = \text{null}$	

Inference Rules

$p = \text{malloc}(\dots)$	$D_kill = \{D[p][s] \mid D[p][s] == 1\} \cup \{D[s][p] \mid D[s][p] == 1\}$ $I_kill = \{I[p][s] \mid I[p][s] == 1\}$
$p = q$	$D_gen = \{D[p][p]\}$ $I_gen = \{I[p][p]\}$
$p = q \rightarrow f$	
$p = \&(q \rightarrow f)$	
$p = q \text{ op } k$	
$p = \text{null}$	$p.shape = Tree$
$p \rightarrow f = q$	
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Inference Rules

$p = \text{malloc}(\dots)$	D_kill and I_kill sets same as for allocation statement.
$p = q$	$D_gen = \{D[s][p] \mid I[s][q] \text{ and } s \neq p\} \cup \{D[p][s] \mid D[q][s] \text{ and } s \neq p \text{ and } s \neq q\} \cup \{D[p][q] \mid q.shape == Cycle\} \cup \{D[p][p] \mid D[q][q]\}$
$p = q \rightarrow f$	$I_gen = \{I[p][s] \mid I[q][s] \text{ and } s \neq p\} \cup \{I[p][p] \mid I[q][q]\}$
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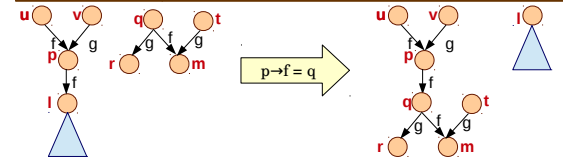


Inference Rules

$p = \text{malloc}(\dots)$	Processing is the same as for $p = q$ statement. This means the analysis loses field-sensitivity.
$p = q$	
$p = q \rightarrow f$	
$p = \&(q \rightarrow f)$	A recent work from IITK (Dasgupta, Karkare, Reddy) addresses this issue.
$p = q \text{ op } k$	
$p = \text{null}$	
$p \rightarrow f = q$	
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Inference Rules

$p = \text{malloc}(\dots)$	$D_kill = \{\}, I_kill = \{\}$
$p = q$	$D_gen = \{D[r][s] \mid D[r][p] \text{ and } D[q][s]\}$
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$p = \&(q \rightarrow f)$	$D[q][p] \text{ and } D[s][q] \Rightarrow s.\text{shape} = \text{Cycle}$
$p = q \text{ op } k$	$D[q][p] \text{ and } D[s][p] \Rightarrow s.\text{shape} = \text{Cycle}$
$p = \text{null}$	$!D[q][p] \text{ and } D[s][p] \text{ and } I[s][q] \text{ and } q.\text{shape} == \text{Tree}$ $\Rightarrow s.\text{shape} = \max(s.\text{shape}, \text{DAG})$
$p \rightarrow f = q$	$!D[q][p] \text{ and } D[s][p] \text{ and } q.\text{shape} != \text{Tree}$ $\Rightarrow s.\text{shape} = \max(s.\text{shape}, q.\text{shape})$
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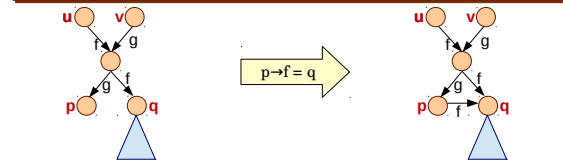


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	Tree	DAG	Cycle
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DAG	DAG	DAG	Cycle
Cycle	Cycle	Cycle	Cycle

$\max(\text{shape1}, \text{shape2})$

Inference Rules

<code>p = malloc(...)</code>	<code>D_kill = { }, I_kill = { }</code>
<code>p = q</code>	<code>D_gen = { }</code>
<code>p = q→f</code>	<code>I_gen = { }</code>
<code>p = &(q→f)</code>	
<code>p = q op k</code>	No changes to the shape of p.
<code>p = null</code>	
<code>p→f = q</code>	
<code>p→f = null</code>	

Example

```
listReverse(List x) {
  assert("x is an acyclic singly linked list");

  for (y = null; x;) {
    t = y;
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  }
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}
```

Interference

x	y	t
x	0	0
y		0
t		

Direction

x	y	t
x	1	0
y	0	0
t	0	0

Shape

x	tree
y	tree
t	tree

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We need to assume a finite representation for the data structure.

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```

Since we do not model fields, we can't say that x is unreachable from y.

Interference

	x	y	t
x		1	0
y			0
t			

Direction

	x	y	t
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x	tree
y	cycle
t	cycle

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Interference

	x	y	t
x		1	1
y			1
t			

Direction

	x	y	t
x	1	1	0
y	1	1	0
t	1	1	1

Shape

x	tree
y	cycle
t	cycle

Example

```
listReverse(List x) {
  assert("x is an acyclic singly linked list");

  for (y = null; x;) {
    t = y;
    y = x;
    x = x->next;
    y->next = t;
  }
  x = y;
  t = null;
  y = null;
}
```

Interference

	x	y	t
x		1	1
y			1
t			

Direction

	x	y	t
x	1	1	1
y	1	1	1
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x	cycle
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Interference

	x	y	t
x		1	1
y			1
t			

Direction

	x	y	t
x	1	0	0
y	1	1	0
t	1	1	1

Shape

x	tree
y	cycle
t	cycle

Improvements

- SSA form
- Field-sensitivity
- Heap modeling

Summary

- Shape analysis helps several transforms.
- Existing techniques often trade off precision for efficiency.
- We are still far away from a precise and scalable analysis.