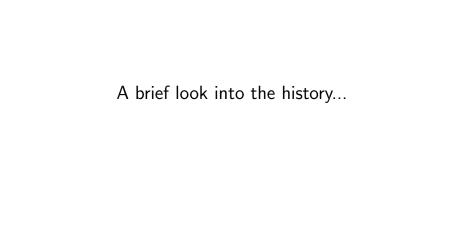
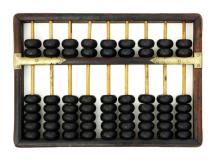
# CS1100 – Introduction to Programming Lecture 2

Instructor: Shweta Agrawal (shweta.a@cse.iitm.ac.in)

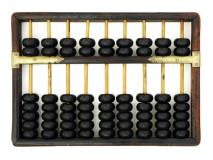


#### From Abacus to Apple



- Counting frame.
- One of the earliest form of calculator.
- Still used by kids to do fast simple arithmetic.

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- Counting frame.
- One of the earliest form of calculator.
- Still used by kids to do fast simple arithmetic.
- Followed by mechanical calculators by B. Pascal (1642),
   G. W. Leibniz (1671).
  - Used cogs / interlocking gears.
  - Performed  $+, -, *, /\sqrt{.}$
  - Leibniz is credited of creating the binary system.

# Jaquard looms (1804)



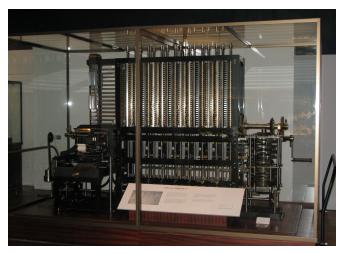


# Charles Babbage (1791–1871)



- Regarded as the "Father of Computer".
- Conceived of a machine that has all the parts of a modern computer, input, a memory, a processor, and an output (1850).

# Difference Engine (1850)



Difference engine built from Babbage's design (London Science Museum).

# Ada Lovlace (1815–1852)



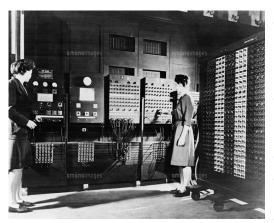
- "Wrote" the description of the mechanical computer of Babbage.
- Regarded as the first programmer ever.
- The programming language ADA is named after her.

# Alan Turing (1912 – 1954)



- Father of Theoretical Computer Science (TCS) and Artificial Intelligence (AI).
- Turing machine a model for a general purpose computer.
- Turing test how intelligent is a machine?

#### First Electronic Computer: ENIAC 1946



Electronic Numerical Integerator

And Calculator.

- 50,000 vacuum tubes, diodes, relays, resistors, capacitors.
- 5 million hand-soldered joints.
- Weighed 27 tons.
- Covered 167 m<sup>2</sup> area.
- Consumed 150 kW of power.

## 1946 - 1976



**Transistors** 



Integrated Circuits

#### 1946 - 1976



**Transistors** 

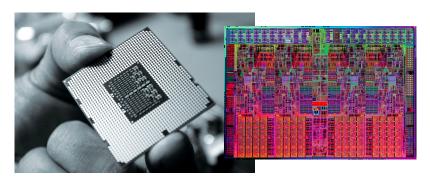


Integrated Circuits



Apple Macintosh

#### Today's World: Core i7 Processor



2008-15: Intel Core i7 Processor

Clock speed: > 2.5 GHz

No. of Transistors: 0.731 - 1.3B

Doubles every two years (Moore's law!)

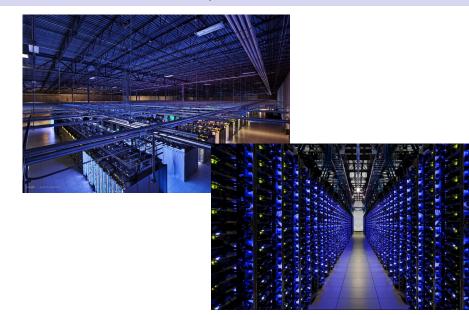
Technology: 45 - 22nm CMOS Area:  $263 - 181mm^2$ .

Nowadays: Multicore (as clock speed increased) with cooling units!

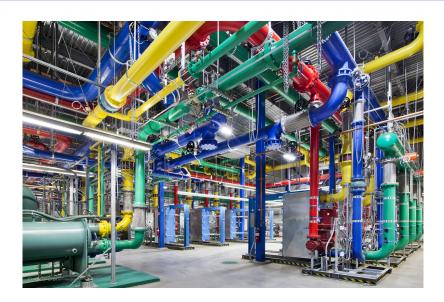
# Modern computing devices



# Data Centers: Processing/Storing Huge volume of data



# Even Cooling them is a big deal ...



# What happens next?



# What happens next?



# How does the computer represent data?

 To store: Numbers, text, graphics and images, video, audio, program instructions.

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- In some way, all information is digitized broken down into pieces and represented as numbers.

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- To store: Numbers, text, graphics and images, video, audio, program instructions.
- In some way, all information is digitized broken down into pieces and represented as numbers.
- Example : Representing Text Digitally.
  - Every character is stored as a number, including spaces, digits, and punctuation.
  - Corresponding upper and lower case letters are separate characters.

Hi, Heather.

#### The ASCII table

American Standard Code for Information Interchange (ASCII).

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Dec	Н	Oct	Cha	r	Dec	Нх	Oct	Html	Chr	Dec	Нх	Oct	Html	Chr	Dec	: Hx	Oct	Html Ch	nr
0	0	000	NUL	(null)	32	20	040	6#32;	Space	64	40	100	e#64;	0	96	60	140	`	
1	1	001	SOH	(start of heading)	33	21	041	6#33;	1	65	41	101	A	A	97	61	141	a	a
2	2	002	STX	(start of text)	34	22	042	6#34;	rr	66	42	102	B	В	98	62	142	b	b
3	3	003	ETX	(end of text)	35	23	043	6#35;	#				6#67;					c	
4	4	004	EOT	(end of transmission)				6#36;		68	44	104	£#68;	D	100	64	144	d	d
5	5	005	ENQ	(enquiry)				6#37;					E					e	
6	6	006	ACK	(acknowledge)				6#38;		70	46	106	6#70;	F				f	
7	7	007	BEL	(bell)				6#39;					6#71;					g	
8	8	010	BS	(backspace)	40	28	050	a#40;	(	72	48	110	6#72;	H	104	68	150	a#104;	h
9	9	011	TAB	(horizontal tab)	41	29	051	6#41;	)	73	49	111	6#73;	I	105	69	151	a#105;	1
10	A	012	LF	(NL line feed, new line)	42	2A	052	6#42;	*	74	4A	112	6#74;	J	106	6A	152	a#106;	j
11	В	013	VT	(vertical tab)	43	2B	053	6#43;	+	75	4B	113	6#75;	K	107	6B	153	a#107;	k
12	C	014	FF	(NP form feed, new page)	44	20	054	6#44;	,	76	4C	114	6#76;	L	108	6C	154	a#108;	1
13	D	015	CR	(carriage return)	45	2D	055	6#45;	-	77	4D	115	6#77;	M	109	6D	155	6#109;	m
14	E	016	50	(shift out)	46	2E	056	a#46;		78	4E	116	6#78;	N	110	6E	156	n	n
15	F	017	SI	(shift in)	47	2F	057	6#47;	1	79	4F	117	6#79;	0	111	6F	157	6#111;	0
16	10	020	DLE	(data link escape)	48	30	060	6#48;	0	80	50	120	£#80;	P	112	70	160	p	p
17	11	021	DC1	(device control 1)	49	31	061	6#49;	1	81	51	121	6#81;	Q	113	71	161	q	q
18	12	022	DC2	(device control 2)	50	32	062	a#50;	2	82	52	122	6#82;	R	114	72	162	a#114;	r
19	13	023	DC3	(device control 3)	51	33	063	3	3	83	53	123	6#83;	S	115	73	163	s	3
20	14	024	DC4	(device control 4)	52	34	064	4	4	84	54	124	T	T	116	74	164	t	t
21	15	025	NAK	(negative acknowledge)	53	35	065	5	5	85	55	125	e#85;	U	117	75	165	u	u
22	16	026	SYN	(synchronous idle)	54	36	066	a#54;	6	86	56	126	6#86;	V	118	76	166	v	V
23	17	027	ETB	(end of trans. block)	55	37	067	6#55;	7	87	57	127	6#87;	W	119	77	167	6#119;	W
24	18	030	CAN	(cancel)	56	38	070	a#56;	8	88	58	130	4#88;	X	120	78	170	x	x
25	19	031	EM	(end of medium)	57	39	071	6#57;	9	89	59	131	4#89;	Y	121	79	171	6#121;	Y
26	1A	032	SUB	(substitute)	58	3A	072	6#58;	:	90	5A	132	6#90;	Z	122	7A	172	6#122;	Z
27	1B	033	ESC	(escape)	59	3B	073	6#59;	;	91	5B	133	6#91;	[	123	7B	173	6#123;	{
28	10	034	FS	(file separator)	60	3C	074	a#60;	<	92	5C	134	6#92;	1	124	7C	174	6#124;	1
29	1D	035	GS	(group separator)	61	3D	075	6#61;	=	93	5D	135	6#93;	]	125	7D	175	6#125;	}
30	1E	036	RS	(record separator)	62	3E	076	6#62;	>	94	5E	136	6#94;		126	7E	176	6#126;	~
31	1F	037	US	(unit separator)	63	3F	077	4#63;	?	95	5F	137	6#95;		127	7F	177	6#127;	DEL

Source: www.LookupTables.com

#### Number Systems.

• Decimal (base 10 - uses 10 symbols  $\{0...9\}$ . Eg : 0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13 ....

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   Eg: 1, 11, 111, 1111, ....

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- Hexadecimal(base 16 uses A-F for 10-15)
   Eg: 0, 1, ..., 9, A, B, C, D, E, F, 10, 11, ... 19, 1A, 1B, ...

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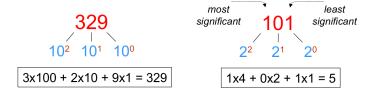
# Quick Primer on Number System : Base n

Take every "digit" and multiply by increasing powers of n and add.

3x100 + 2x10 + 9x1 = 329

## Quick Primer on Number System : Base n

Take every "digit" and multiply by increasing powers of n and add.



## Converting from Decimal to Binary

Conver the decimal number 39 to binary (base 2).

2 
$$39$$
  $39 = 2*19 + 1$   $= 2*(2*9 + 1) + 1$   $= 2^2*9 + 2^1*1 + 1$   $= 2^2*(2*4 + 1) + 2^1*1 + 1$   $= 2^3*4 + 2^2*1 + 2^1*1 + 1$   $= 2^3*(2*2 + 0) + 2^2*1 + 2^1*1 + 1$   $= 2^4*2 + 2^3*0 + 2^2*1 + 2^1*1 + 1$   $= 2^4*(2*1 + 0) + \dots$   $= 2^5*1 + 2^4*0 + 2^3*0 + 2^2*1 + 2^1*1 + 1$   $= (1 \times 2^0) + (1 \times 2^1) + (1 \times 2^2) + (0 \times 2^3) + (0 \times 2^4) + (1 \times 2^5)$ 

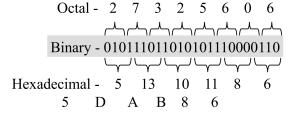
$$(100111)_2 = (1 \times 2^0) + (1 \times 2^1) + (1 \times 2^2) + (0 \times 2^3) + (0 \times 2^4) + (1 \times 2^5)$$
  
=  $(39)_{10}$ 

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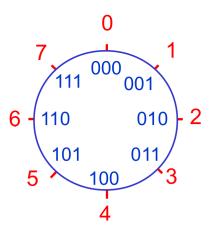


# Representing values in Binary

If we have m bits, we can represent  $2^m$  unique different values.

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If we have m bits, we can represent  $2^m$  unique different values. A useful circle :



### Sign Magnitude notation

• Use one bit for sign, others for magnitude of the number.

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		Sign Magn.
0	0	0
0	1	+1
1	0	+2
1	1	+3
0	0	0
0	1	-1
1	0	-2
1	1	-3
	0 1 1 0	0 1 1 0 1 1 0 0 0 1

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• Use one bit for sign, others for magnitude of the number.

			Sign Magn.
0	0	0	0
0	0	1	+1
0	1	0	+2
0	1	1	+3
1	0	0	0
1	0	1	-1
1	1	0	-2
1	1	1	-3

- using *n* bits:  $-(2^{n-1}-1)...(2^{n-1}-1)$ .
- zero has two representations.

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• for a negative number n, represent the number by the bit complement of its binary rep. using k bits.

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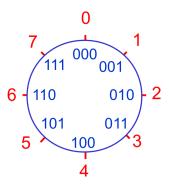
			Sign Magn.	Ones comp.
0	0	0	0	0
0	0	1	+1	+1
0	1	0	+2	+2
0	1	1	+3	+3
1	0	0	0	-3
1	0	1	-1	-2
1	1	0	-2	-1
1	1	1	-3	0

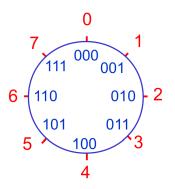
#### Ones complement notation

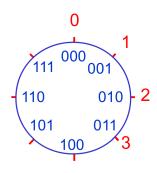
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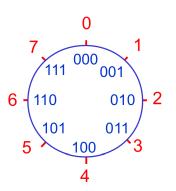
			Sign Magn.	Ones comp.
0	0	0	0	0
0	0	1	+1	+1
0	1	0	+2	+2
0	1	1	+3	+3
1	0	0	0	-3
1	0	1	-1	-2
1	1	0	-2	-1
1	1	1	-3	0

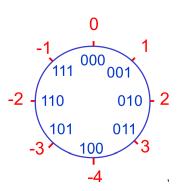
- using *n* bits:  $-(2^{n-1}-1)\dots(2^{n-1}-1)$ .
- zero has two representations.
- not very widely used representation.











#### Twos complement notation

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- for a negative number -n, represent the number as  $2^k n$ .

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			Sign Magn.	Ones comp.	Twos comp.
0	0	0	0	0	0
0	0	1	+1	+1	+1
0	1	0	+2	+2	+2
0	1	1	+3	+3	+3
1	0	0	0	-3	-4
1	0	1	-1	-2	-3
1	1	0	-2	-1	-2
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1	0	0	0	-3	-4
1	0	1	-1	-2	-3
1	1	0	-2	-1	-2
1	1	1	-3	0	-1

- using *n* bits:  $-(2^{n-1}) \dots (2^{n-1}-1)$ .
- widely used representation.

### Arithmetic with these representations

			Sign Magn.	Ones comp.	Twos comp.
0	0	0	0	0	0
0	0	1	+1	+1	+1
0	1	0	+2	+2	+2
0	1	1	+3	+3	+3
1	0	0	0	-3	-4
1	0	1	-1	-2	-3
1	1	0	-2	-1	-2
1	1	1	-3	0	-1

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			Sign Magn.	Ones comp.	Twos comp.
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0	1	1	+3	+3	+3
1	0	0	0	-3	-4
1	0	1	-1	-2	-3
1	1	0	-2	-1	-2
1	1	1	-3	0	-1

• 
$$2 + (-3)$$

### Arithmetic with these representations

			Sign Magn.	Ones comp.	Twos comp.
0	0	0	0	0	0
0	0	1	+1	+1	+1
0	1	0	+2	+2	+2
0	1	1	+3	+3	+3
1	0	0	0	-3	-4
1	0	1	-1	-2	-3
1	1	0	-2	-1	-2
1	1	1	-3	0	-1

- 2 + (-3)
- 3 + (-2)

# More examples : The case of 4 bits

1 1 √ √ √	corresp. dec. oper.		esp. oper.
0110	+6	0100	+4
+1101	+ -3	+1001	+ -7
10011 =	+3 +3	1101 = -3	-3
cor	rect result	correct res	sult
Exa	ample (a)	Example (	<b>b</b> )

# More examples : The case of 4 bits

	corresp. dec. oper.		orresp. ec. oper.
0011	+3	1110	-2
+0100	+ +4	+1010	+ -6
0111 = +	7 +7	11000 = -8	-8
corre	ct result	correct i	esult
Exam	ple (c)	Example	e (d)

# More examples: The case of 4 bits

$$\frac{1}{\sqrt{1}}$$
 Corresp. dec. oper.

  $\frac{1101}{1010}$ 
 $\frac{-3}{1010}$ 
 $\frac{+5}{1010}$ 
 $\frac{10111}{1011}$ 
 $\frac{+6}{10110}$ 
 $\frac{+6}{10111}$ 

 incorrect result
 incorrect result

 Example (e)
 Example (f)

**Overflow Detection Rule**: If two numbers with the same sign (both positive or both negative) are added, then overflow occurs if and only if the binary representation of the result has the opposite sign.

### What to do?

**How to Detect it?** : The technique of overflow detection is easily implemented in electronic circuitry, and it is a standard feature in digital adder circuits.

How to Prevent it?: Use more bits!